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A 2.56 GBPS SERIAL WIRELINE TRANSCEIVER THAT  
SUPPORTS AN AUXILIARY CHANNEL  
AND  
A HYBRID LINE DRIVER TO COMPENSATE LARGE CHANNEL LOSS

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A 2.56 GBPS SERIAL WIRELINE TRANSCEIVER THAT  
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A Dissertation Presented to the Graduate Faculty of

Lyle School of Engineering

Southern Methodist University

in

Partial Fulfillment of the Requirements

for the degree of

Doctor of Philosophy

with a

Major in Electrical Engineering

by

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August 4, 2020

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Xiaoran Wang

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It has been five years since I started to pursue my PhD degree. During the past five years, my life has changed a lot. At the same time, the research work also opens a new world for me, which is full of unknown, challenges, excitements and hopes. When I look back, there are always the people helping me a lot.

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## STATEMENT BY AUTHOR

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B.S., South China University of Technology, 2013

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A 2.56 Gbps Serial Wireline Transceiver that Supports An Auxiliary Channel  
and  
A Hybrid Line Driver to Compensate Large Channel Loss

Advisor: Professor Ping Gui

Doctor of Philosophy conferred August 4<sup>th</sup>, 2020

Dissertation completed August 4<sup>th</sup>, 2020

Serial transceiver links are widely used for high-speed point-to-point communications. This dissertation describes two transceiver link designs for two different applications.

In serial wireline communications, security is an increasingly important factor of concern. Securing an information processing system at the application and system software layers is regarded as a necessary but incomplete defense against the cyber security threats. Many security measures at the hardware level require an additional data channel with extra bandwidth to transmit authentication information or to just increase redundancy. A competing issue is that most existing designs cannot support the additional channel required to implement these measures without an expensive redesign effort.

In this dissertation, an asynchronous serial transceiver that is capable of transmitting and receiving an auxiliary data stream concurrently with the primary data stream is described. The transceiver instantiates the auxiliary data stream by modulating the phase of the primary data without affecting the primary channel transmission and recovery mechanisms. Standard receiver interoperability is maintained since the auxiliary data appears as primary data jitter. Analysis of



the proposed transceiver and considerations of the system parameters are included and can be used to determine how such an auxiliary channel is implemented. The proposed transceiver with the auxiliary channel can be widely used in many data communication applications such as for transmitting signatures for authentication or other control information, steganography, or additional data in an existing serial link. A prototype transceiver, implemented in a 65 nm CMOS process, demonstrates the proposed concept with an 80 Mbps auxiliary channel in a 2.56 Gbps asynchronous serial link.

The Deep Underground Neutrino Experiment (DUNE) requires that the front-end transmitters operate at cryogenic temperature and drive 25-35 meters long twin-axial (twinax) cables. To compensate the frequency-dependent channel loss over the long cables and alleviate the de-emphasizing of the low-frequency signal magnitude, a hybrid of a current-mode (CM) transmitter equalization (TXEQ) and a voltage-mode (VM) pre-emphasis is proposed. The TXEQ employs a finite-impulse response (FIR) filter to boost the high-frequency components while de-emphasizing the low-frequency signal magnitude, thereby flattening the overall channel frequency response and reducing the Intersymbol Interference (ISI). The VM pre-emphasis is proposed to further mitigate ISI by boosting the high-frequency portion without degrading the signal magnitude, allowing for high signal swing. The main driver utilizes VM source-series-terminated (SST) output stages, which offers higher signal swing and better power efficiency than the conventional current-mode logic (CML) drivers. To ensure the lifetime and reliability at cryogenic temperature, the transmitter is implemented in a 65-nm CMOS process operating at 1.1 V of supply voltage and employing transistors with larger than minimum lengths. Silicon measurement results have validated the proposed approaches.

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This is dedicated to my parents.



## CHAPTER 1 INTRODUCTION

### 1.1 Motivation

Serial wireline communications are widely used in most of the computer networks as well as the chip-to-chip communications. In serial wireline communications, in addition to data rate and data transfer quality (Bit Error Rate), security is another increasingly important factor of concern. Securing an information processing system at the application and system software layers is regarded as a necessary but incomplete defense against the cyber security threats. Encryption is a commonly employed method with the goal of preventing unauthorized access to sensitive information [1]-[2]. However, the modification or redesign of an existing system to include encryption at the hardware layer can add significant expense and result in compatibility issues with other systems and specifications as well as interoperability issues with other contemporary versions of similar systems. Moreover, sometimes the mere presence of an encrypted channel provides an adversary with information that is undesired and encourages increased attacks [3]-[4]. More emphasis is being placed in the area of hardware security due to the emergence of exploits at these lower layers of data transmission and processing [5]-[9].

Many security measures at the hardware level require an additional data channel with extra bandwidth to transmit authentication information or to just increase redundancy. A competing issue is that most existing designs cannot support the additional channel required to implement these measures without an expensive redesign effort. Even if such redesign efforts are

accomplished, the resulting more secure products may not be backwardly compatible with earlier or standard generations.

Modern information processing circuitry is becoming very common for information exchanges to be accomplished via asynchronous serial links [10]-[14]. Asynchronous serial transceivers are ubiquitous in today's devices and are used to interface between blocks within and between integrated circuits (ICs), and between packaged systems, typically using industry standards such as USB, MIPI, and PCI-e [15].

The motivation of the first work of this dissertation is to create a new approach and transceiver architecture to address the above challenges and problems, and enhance the security of serial wireline communications and to ensure authenticated data transfer at the circuit level. The proposed new serial transmitter is able to embed a signature in the serial data stream by modulating the last-stage high-speed clock in the serializer, and the proposed receiver can recover not only the transmitted data and clock but also the embedded signature. The modulation of the clock signal results in higher deterministic jitter on the transmitter side. This imposes challenges in the transmitter and receiver design. On the transmitter side, it needs to produce small clock phase modulation which can increase the deterministic jitter in the transceiver. On the receiver side, it needs to recover the actual transmitted data in the presence of higher jitter compared to the case without clock phase modulation, and at the same time detect the clock phase modulation as the signature.

The output driver of a data link transmitter is a critical part that determines the overall performance of the entire high-speed serial link. Many applications require the transmitter driver to consume low power and operate at high speed. In addition, most of the applications require the

transmitter driver to provide equalization to compensate the frequency-dependent channel loss inherent with the transmission channels thus mitigating intersymbol interference (ISI). Meanwhile, large signal swing from the transmitter is also desirable in order for the receiver to recover the data stream with a good bit-error rate (BER) performance.

The Deep Underground Neutrino Experiment (DUNE) is a dual-site experiment consisting of two sets of detectors: one in Illinois (near detectors) and the other one in South Dakota (far detectors). The far detectors will be in a single-phase time projection chamber, with 10K tons of liquid argon operating at cryogenic temperature (89 Kelvin). The interaction of the neutrino particles in liquid argon is detected and converted to digital data streams inside the chamber. At the transmitter, the digital streams are combined to a 1.28 Gbps serial data stream and sent to the warm interface of the chamber cryostat. Due to the large dimensions of the chamber cryostat, the line drivers in the readout ICs need to drive 25-35 meters long twin-axial (twinax) cables to transmit the high-speed data streams to the Warm Interface Boards (WIB).

The motivation of the second work of this dissertation is to design a line driver for DUNE. Compensating large channel loss while maintaining large signal swing (thus ensuring good eye diagrams) are required for the transmitters used in DUNE. In addition, while major portion of the cables will be immersed into liquid argon at cryogenic temperature (89 Kelvin), a small portion of the cables may be at the warm side with uncertain temperature and to be connected to the WIB. Given the high and uncertain insertion loss resulting from a large range of the cable lengths and a complex temperature environment, reconfigurable equalization at the transmitter is needed in DUNE.

## 1.2 Research Contribution

This dissertation proposed a new asynchronous serial transceiver that is capable of transmitting and receiving an auxiliary data stream along with the primary data stream on a single asynchronous serial link. The proposed transceiver embeds the auxiliary data stream by modulating the phase of the primary data in accordance with it. The receiver recovers both the primary and auxiliary data simultaneously. In a standard receiver, which is not equipped with the phase demodulation capability, the auxiliary data appears as jitter of the primary data. The jitter caused by the auxiliary data still falls within the jitter budget of the transceiver, and having this much jitter would not adversely affect the functionality of the primary data recovery. The proposed system can be widely used in many data communication applications such as for transmitting a hidden signature for data authentication, or as control and/or additional data in an existing serial link. A prototype transceiver, implemented in a 65 nm CMOS process, demonstrates the proposed concept with 2.56 Gbps primary data and 80 Mbps auxiliary data channels.

This dissertation also proposed a hybrid line driver to drive at least 25-meter long twin-axial (twinax) cables in Deep Underground Neutrino Experiment. To compensate the high-frequency loss over the long cables and alleviate the de-emphasis of the low-frequency signal magnitude, a hybrid of current-mode transmitter equalization (TXEQ) and voltage-mode pre-emphasis is proposed. The TXEQ employs a finite-impulse response (FIR) filter to boost the high-frequency components while deemphasizing the low-frequency signal magnitude, thereby flattening the overall channel frequency response and reducing Inter-Symbol Interference (ISI). Voltage-mode pre-emphasis is proposed to further boost the high-frequency portion without degrading the signal magnitude. The main driver utilizes voltage-mode source-series-terminated (SST) output stages, which offers higher signal swing and better power efficiency than the conventional current-mode

logic (CML) drivers. Designed in 65-nm CMOS, the proposed hybrid line driver operates at 1.28 Gbps and with 4.1 mW power consumption when driving a 25-meter long twinax cable at cryogenic temperature.

### **1.3 Dissertation Organization**

This dissertation is organized as follows. Chapter 1 is the introduction, including the motivation, research contribution and organization of the dissertation. Chapter 2 presents a review of the serial wireline communication system. The key parts of the serial wireline communication system are also introduced, such as serializer/deserializer (SerDes), clock and data recovery (CDR) and the line driver. In chapter 3, the 2.56 Gbps serial wireline transceiver that supports an auxiliary channel is shown, and the measurement results are demonstrated. Chapter 4 depicts the hybrid line driver to compensate large channel loss. In chapter 5, the summary and future work of the dissertation are presented.

## CHAPTER 2 SERIAL WIRELINE COMMUNICATION SYSTEM

### 2.1 SerDes

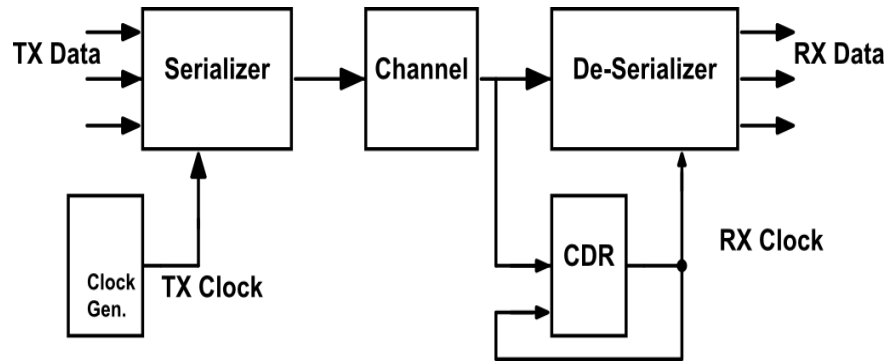


Figure 1. Block diagram of serial data transmission system.

Figure 1 shows a generic block diagram of a high-speed wireline data transmission system. The left side is the transmitter and the right side is the receiver. Those two parts constitute the serial wireline communication system. At the transmitter (TX) side, Parallel digital signals are converted to a binary sequence through a data serializer for the single-channel serial-data transmission. The data to be sent is bundled into a high-speed stream in the transmitter, and then transmitted to the channel by the driver circuit. At the receiver (RX) side, the CDR recovers the clock signal and retimes the data.

SerDes is used in the transmitter and receiver of the high-speed serial wireline data communication system. The primary use of the SerDes is to provide data transmission over a single channel in order to minimize the number of I/O pins and interconnects.

The input of the serializer is an  $n$ -bit data path that is serialized to a one-bit serial data signal. Generally, the value of  $n$  is a multiple of 8 or 10, and may be programmable on some implementations [18]. Values of  $n$  that are multiples of 8 are useful for sending unencoded data bytes; values of  $n$  that are multiples of 10 are useful for protocols using 8B/10B coding. In the serializer, shift registers receive the parallel input data and clock, and then shift the data out to a higher serial clock frequency. As shown in Figure 2, the parallel input data is assumed to be 4 bits, and the input reference clock rate is 500MHz. In order to transfer the parallel data into serial data, the first set of two shift registers use 1GHz clock and transfer the data to 2-bit parallel data, and the second set of shift register uses 2GHz clock and transfers the data to 1-bit serial data.

The deserializer first extracts the clock from the incoming serial data stream using a CDR and uses this clock to deserializer the data into parallel bits at lower frequency. The CDR working principles will be introduced in the following chapters.

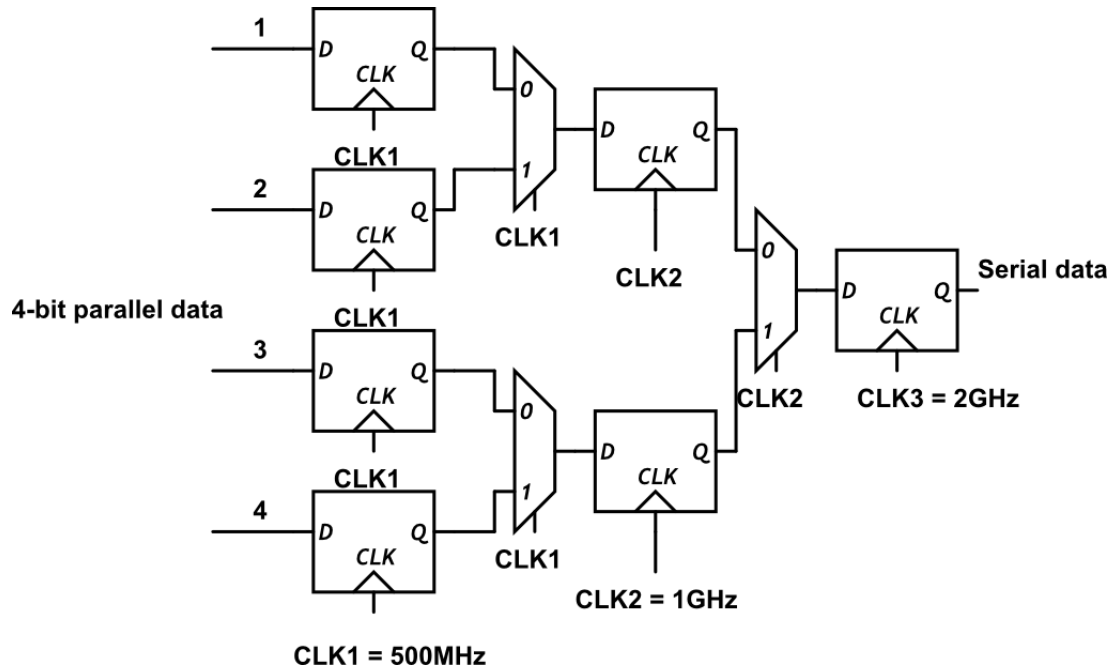


Figure 2. Block diagram of serializer

## 2.2 Clock and Data Recovery

In the high-speed wireline communication systems, such as wireline long-haul networks, and chip-to-chip or backplane communications, the received data are distortion [19] [20] [21]. The data is distorted because of many factors such as noises of devices, limited bandwidth of channels, signal reflection, duty-cycle distortion, cross-talk, and power supply noise. A clock and data recovery (CDR) circuit is a critical block in the data link to extract the timing information and retime the data so as to recover the distorted data.



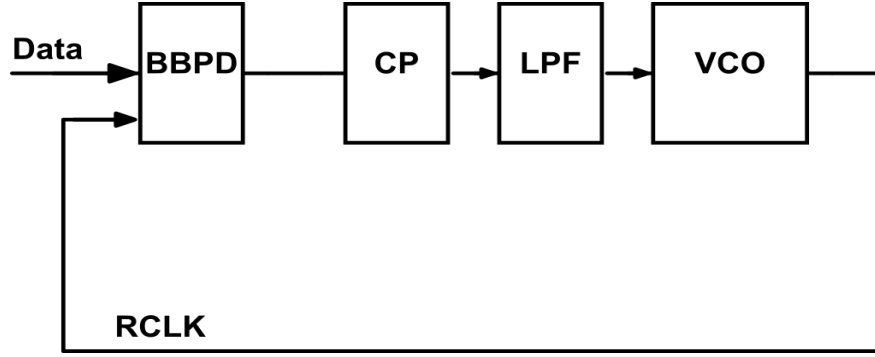


Figure 3. PLL-based CDR structure

PLL (Phase locked loop) - based CDR is a topology using feedback phase tracking. Figure 3 shows a PLL-based CDR that is a single negative feedback loop without extra PLL to generate the clock signal. PLL-based CDR is the most common scheme of CDR based on a second-order PLL circuit. The tracking loop employs a phase detector (PD), a charge pump (CP) and a low pass filter (LPF) to drive the voltage control oscillator (VCO) frequency towards the input data rate. When the phase error falls into the capture range of phase tracking loop, the phase tracking loop aligns the phase of clock with the data. In PLL-based CDR, the output of the bang-bang phase detector makes the charge pump produce ‘up’ or ‘down’ signal, and those ‘up’ and ‘down’ signal will tune the VCO frequency and adjust the feedback clock signal [22], [23], [24].

PLL-based CDR consumes large layout area. Compared to other types of CDR, PLL-based CDR is suitable for single channel CDR circuit and consumes less power.

The jitter performance of PLL-based CDR is worse than other types of CDR architecture. In PLL-based CDR, the stabilizing zero in the forward path introduces jitter peaking and the corner frequencies of the jitter transfer and jitter tolerance transfer functions are at the same frequency, which is the second pole of the PLL. While the bandwidth requirements for good jitter transfer and good jitter tolerance are opposite. Good jitter transfer function needs narrow bandwidth and good

jitter tolerance function needs wide bandwidth, so optimizing one will degrade the other one.

## 2.3 Jitter

Jitter is the time deviation caused by the system noise or other system variation, and is the difference of a signal's ideal or expected arrival time with real arrival time. As can be seen in Figure 4, jitter in a clock signal represents the deviation of the zero crossings from their ideal position in time.

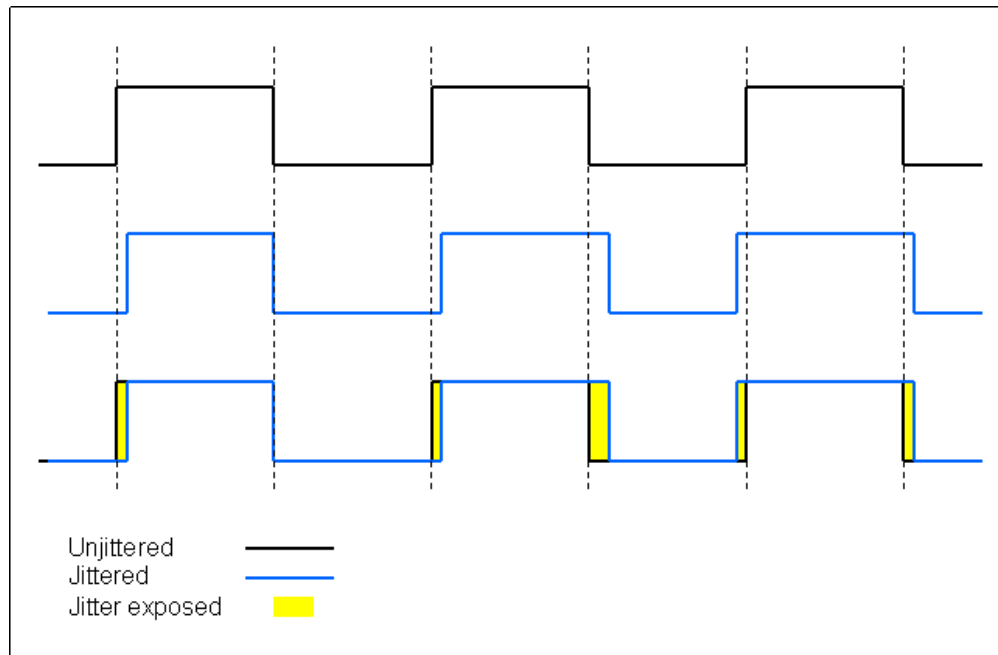


Figure 4. Comparison of the pulse signals with and without jitter

There are two types of jitter, one is random jitter, and the other is deterministic jitter. Random jitter (RJ) fulfills the Gaussian distribution, and is measured by RMS value. Deterministic jitter (DJ) is always fixed in a constant value, and can be quantified by the peak to peak value. The sum of DJ and RJ is total jitter (TJ). TJ is also quantified by the peak to peak value, and is related to the data link BER design spec.

The relationship of TJ, RJ and DJ can be expressed by the following equation.

$$TJ_{pp} = DJ_{pp} + Q_{ber} * RJ_{rms} \quad (1.1)$$

Where  $TJ_{pp}$  is the peak to peak value of the total jitter.  $DJ_{pp}$  is the peak to peak value of the deterministic jitter.  $RJ_{rms}$  is the RMS value of the random jitter.  $Q_{ber}$  is the amount of eye closure due to random jitter that must be accounted for a target bit error rate (BER). A common BER in wireline communication standards is  $10^{-12}$ . The corresponding  $Q_{BER}$  is 14.

The CDR is to recover the data that contains the jitter from the system, and the CDR itself can also produce some jitter, such as feedback clock jitter. System and CDR jitter performance can influence the target BER spec. CDR jitter performance can be categorized into jitter transfer (JTRAN), jitter generation (JGEN) and jitter tolerance (JTOL).

### 2.3.1 JTRAN

JTRAN is the jitter transfer function of the CDR, which is the ratio of output to input jitter as a function of frequency. The JTRAN of the CDR acts as a low pass effect that only the slow jitter in the data can pass through the CDR loop. On the other hand, high frequency jitter is filtered and may cause sampling error if its amplitude is large enough.

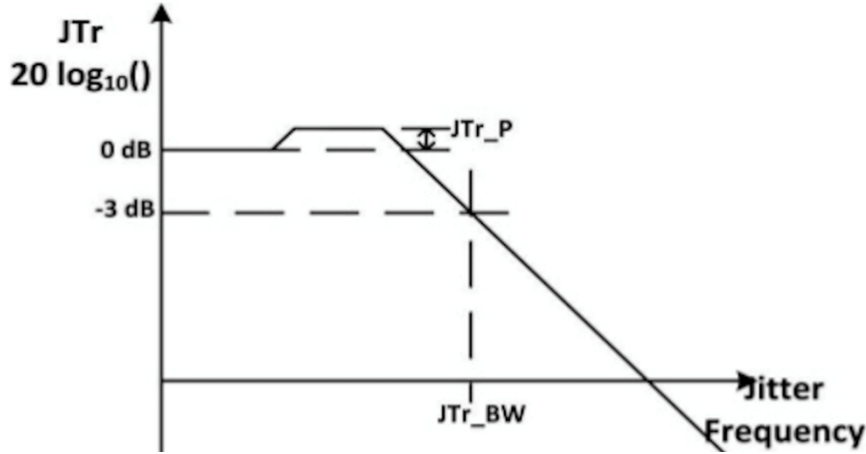


Figure 5. Jitter transfer function

The jitter transfer function of a CDR is shown in Figure 5. As can be seen, CDR loop can track the low speed of signal change. Therefore, the phase of input and the phase of the feedback are the same, and the ratio of the output phase over the input phase at low frequency is 1. As the jitter frequency increases and becomes higher than the CDR loop bandwidth, the CDR loop cannot respond fast enough to the input, and then the ratio of the feedback phase over the input phase turns small. Therefore, the JTRAN starts to decrease. For the CDR design, the JTRAN is a key parameter that defines the jitter bandwidth.

### 2.3.2 JGEN

Jitter comes from everywhere in the communication system, also in CDR circuit itself. JGEN is defined as the jitter generated by CDR itself. As the jitter from other place, CDR jitter is also composed of random jitter (RJ) and deterministic jitter (DJ). When the bang-bang phase detector is used as the phase detector (PD) in the CDR loop, its limit cycle oscillation can produce the deterministic jitter. The random jitter comes from the charge pump, the thermal noise of the low pass filter, as well as the noise of VCO.

### 2.3.3 JTOL

JTOL indicates the jitter tolerance performance. JTOL can be measured by adding the sinusoidal jitter signal of various magnitudes and frequencies to the data, and need to fulfill the target BER requirement at the same time. The jitter tolerance mask is the standard to require the system JTOL performance.

If the CDR jitter transfer function is available, the JTOL can be derived as follows. In order to make sure there is no sampling phase error, the phase different between the input and feedback signal should be less than 0.5UI (unit interval).

$$\Phi_{in} - \Phi_{out} \leq 0.5UI \quad (1.2)$$

If the jitter transfer function  $A(f)$  is already known, the  $\Phi_{out}$  can be replaced by  $\Phi_{in} A(f)$ .

$$\Phi_{in}(1 - A(f)) \leq 0.5UI \quad (1.3)$$

So, the input phase requirement can be derived below,

$$\Phi_{in} \leq \frac{0.5UI}{1 - A(f)} \quad (1.4)$$

The limitation of  $\Phi_{in}$ ' value is the CDR JTOL. For the jitter transfer function we talked above, the low frequency of the jitter,  $A(f)=1$ , and  $\Phi_{in}$  could be infinity. When the jitter frequency exceed the bandwidth of this jitter transfer function, the value of  $A(f)$  start to decrease, and  $\Phi_{in}$  also start to decrease.

Jitter mask is the specification of the system jitter. Figure 6 shows system jitter tolerance mask. JTOL value at any frequency needs to be larger than the jitter mask value.

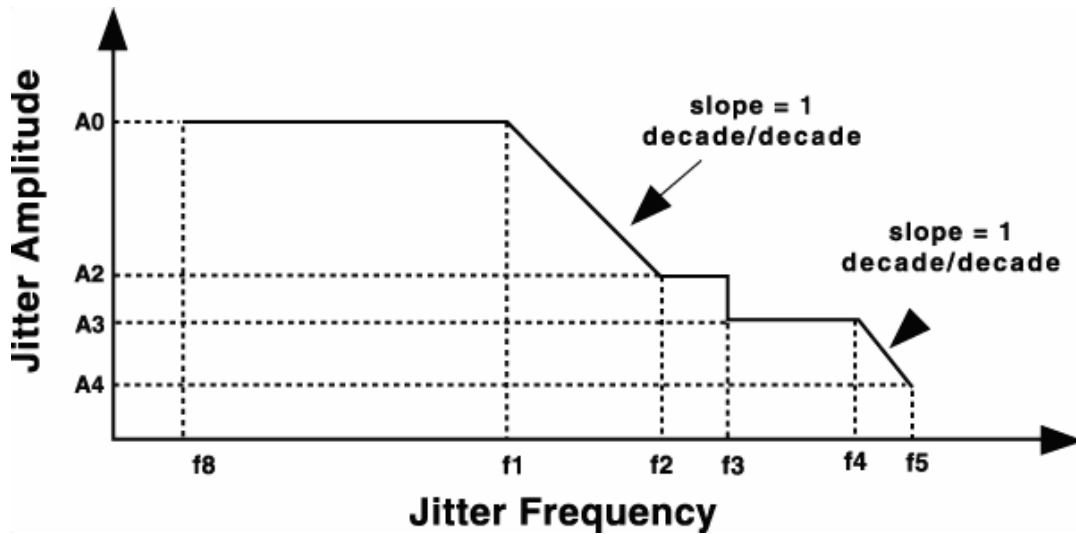


Figure 6. System jitter tolerance mask

## 2.4 Line Driver

The output line driver of a data link transmitter is a critical part that determines the overall performance of the entire high-speed serial link. Not only the drivers need to provide FIR equalization to compensate the frequency-dependent channel loss inherent with the transmission channels thus mitigating ISI, many applications require the transmitter drivers to provide large signal swing for the receiver to recover the data streams with good bit-error rate (BER) performance. Meanwhile, low power consumption and high-speed operation are also desirable in transmitter.

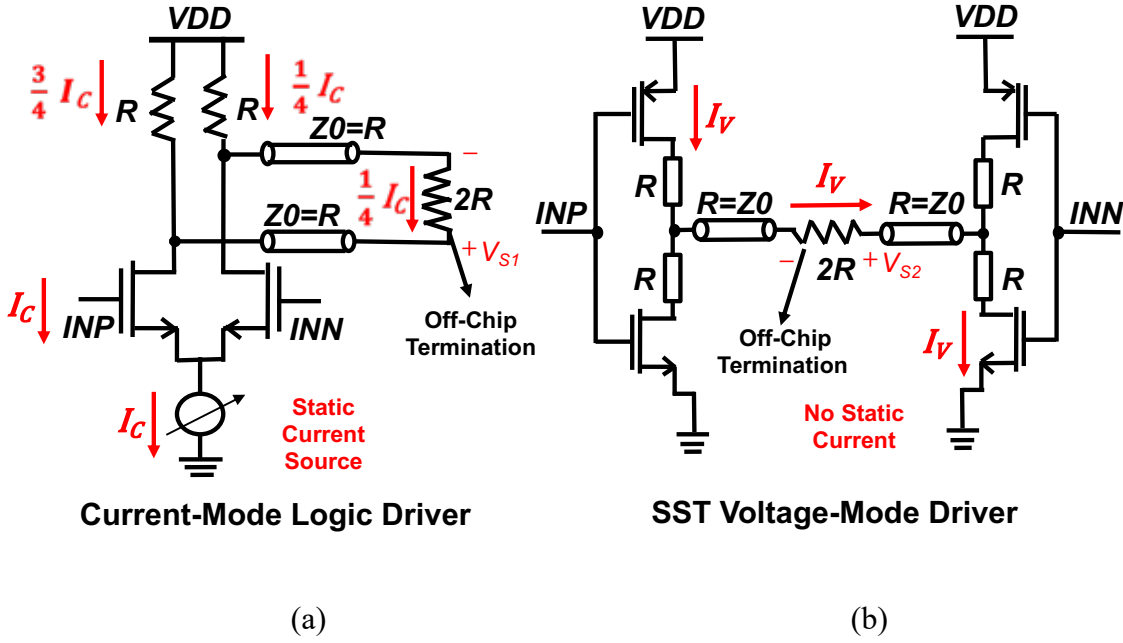


Figure 7. Schematic of (a) CML driver and (b) SST VM driver.

The current-mode (CM) and voltage-mode (VM) drivers are two common design styles for transmitter output drivers [34]-[44]. Conventional high-speed drivers design typically utilizes current-mode-logic (CML) to implement both the main driver and the TXEQ. A simple CML driver schematic is shown in Fig. 7 (a), where a current source provides constant and static current  $I_C$  for the switching transistors. The resistor  $R$  acts as the output impedance of the driver, and its value can be constant and not affected by the value of the static current that determines the tap coefficients in TXEQ.  $Z_0$  is the characteristic impedance for each channel, which is typically equal to the driver's output impedance  $R$ . The off-chip differential termination resistor is  $2R$ . The output swing of the driver is determined by the static current and the value of  $R$ , and the common-mode voltage is determined by the supply voltage as well. The DC current flow  $I_C$  of the CML driver is labelled in Fig. 7 (a), corresponding to the case when the switching transistor on the left branch is fully turned on, and the one on the right branch is turned off. The single-ended signal swing  $V_{S1}$

across the off-chip termination is  $\frac{I_C R}{2}$ . The differential signal swing  $V_{DI}$  is twice of  $V_{SI}$ , thus  $I_C = \frac{V_{DI}}{R}$ .

On the other hand, the VM drivers use the power supply as a voltage source without a constant tail current source, which improves power efficiency since there is no static current being consumed [36]-[40]. An example of a differential VM source-series-terminated (SST) driver is shown in Fig. 7 (b). Each branch of the SST output stage comprises a pull-up PMOS and a pull-down NMOS switching transistor. Both the pull-up and pull-down paths are connected by a series termination resistor  $R$ . The output impedance of the VM driver is the overall resistance of the parallelly connected  $R$  from each driving cell. The switching transistors in the SST output stages should be designed with large size to make their resistance negligible compared to the series resistor  $R$ . When  $INP$  is logic '1' and  $INN$  is logic '0', the DC current flow  $I_V$  is labelled in Fig. 7 (b). The single-ended signal swing  $V_{S2}$  across the off-chip termination resistor is  $2I_V R$ , thus  $I_V = \frac{V_{D2}}{4R}$ , where  $V_{D2}$  is the differential signal swing of the SST VM driver. As can be seen from the above equations, the VM SST driver can be ideally four times power efficient than CML driver when achieving the same signal swing ( $V_{DI}=V_{D2}$ ). In addition, the CMOS-oriented SST driver is flexible to support different termination voltages and able to offer an output swing up to rail-to-rail differentially [41]-[43].



## CHAPTER 3    A SERIAL WIRELINE TRANSCEIVER THAT SUPPORTS AN AUXILIARY CHANNEL

### 3.1 Introduction

To meet the demand of using the channel bandwidth to create an auxiliary channel, we devised and implemented a wireline transceiver for an asynchronous serial channel that provides additional data bandwidth through inclusion of an auxiliary data channel and is interoperable with non-equipped transceivers in earlier generation systems. An auxiliary channel at a lower layer generally provides increased bandwidth to support requirements of security and other system modifications. This technique is also a means for steganography since it allows for communications to be hidden whether they are actually encrypted or not [8]-[9].

Serial communication channels must have a bandwidth in excess of the signal bandwidth that they transmit to allow for reliable communications in accordance with Shannon's capacity theorem. Because signals are always transmitted in the presence of noise, practical channels are designed with a bandwidth margin to account for reliable detection of the transmitted bitstream. It is desirable to efficiently utilize the available bandwidth in asynchronous serial channels. The common trend of the modern wireline transceiver is to be designed with increasingly large data transmission bandwidth on the serial channel. However, besides increasing the transmitted data rate on a single channel, we consider that the available bandwidth can also be used to support an auxiliary channel that is capable of transmitting and receiving an auxiliary data stream in parallel with the primary channel on a single serial link.

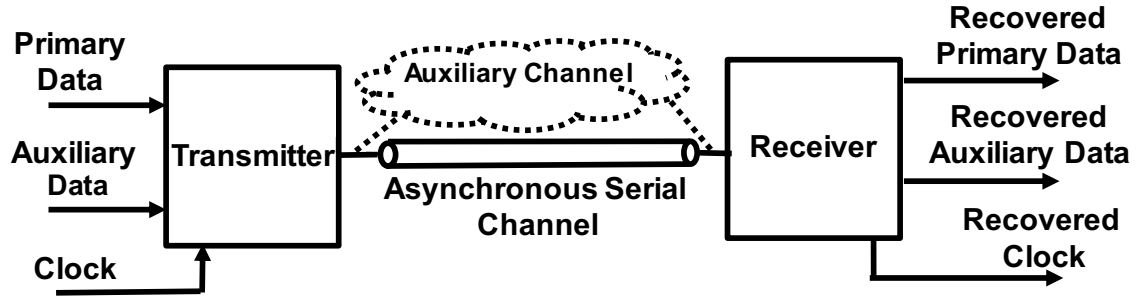


Figure 8. Proposed asynchronous serial link with auxiliary channel.

As shown in Fig. 8, the proposed transceiver system transmits the primary and auxiliary data stream through a single asynchronous serial channel and simultaneously recovers both of them at the receiver.

The proposed novel transmission scheme can be used to provide several benefits mentioned as follows. For example, it can be applied in different ways to enhance hardware security, such as by using the auxiliary channel to carry authentication information for the simultaneously received primary data [16]. Other enhancements may exploit the fact that the auxiliary data transmission can also be considered as a form of steganography since the proposed method provides backward-compatibility with standard transceivers for the primary data channel, and the auxiliary data appears as jitter to a non-equipped transceiver [8]-[9]. Alternatively, the auxiliary channel may be used for error correction or detection on the primary channel, thereby enhancing its data integrity

and capacity without an increase in bandwidth [17]. Moreover, data throughput can also be increased through the use of the additional auxiliary channel.

There are many applications for an auxiliary channel over a serial wireline link. As an example, various schemes involving data authentication at the physical level are desirable without causing significant cost increases. Data authentication can be accomplished via the receiver verifying the data source authenticity through signatures or other data being transmitted over the auxiliary channel. One example is the increasingly common use of reconfigurable FPGAs wherein configuration bitstreams are dynamically loaded from on-board serial memory. Such bitstreams could be accompanied with authenticating signatures that provide some level of security to the device being configured. If the auxiliary data stream comprises sensitive information, it can be further secured through the use of encryption.

Auxiliary channel data transmissions may also be used for purposes other than authentication, for example link quality indices (LQI) or error detection and correction checksums could be transmitted to accompany the primary data [4].

An overall advantage of embedding a physical-layer auxiliary channels in an asynchronous serial link is the use of the channel to transmit additional data in an existing design without the need for including another primary channel. Since transistors are relatively plentiful as compared to on- and off-chip communication channels, it is advantageous to use existing communications channels rather than incorporating costly additional links. As an example, high-speed data on the primary channel may be accompanied with lower bandwidth control or synchronizing data on the auxiliary channel. In this latter case, keeping the auxiliary transmissions hidden may not be a

primary goal, rather the decrease in cost and more efficient usage of available bandwidth could be the motivating factor.

The primary contributions of this work include the design, simulation, analysis and prototype measurement of a new architecture for an asynchronous serial transceiver that supports an auxiliary channel. The potential applications to support hardware-level security is described in the manuscript. We also emphasize that this additional channel is provided in a way that offers backward compatibility, interoperability with non-equipped designs, and minimal redesign of existing systems. The architectures of the transmitter and receiver is described in detail. The guidelines to select the system parameters on the proposed scheme is provided. The analysis on the limits and impact of the additional auxiliary data channel is given. The measurement results demonstrate the function and performance of this proposed transceiver.

### **3.2 Channel Bandwidth Margin**

The Shannon-Hartley channel capacity theorem  $C=B \cdot \log_2(1+SNR)$ , describes the relationship between the bandwidth of a channel versus the signal-to-noise ratio (SNR) of the transmitted signal, where  $C$  represents the channel capacity in bit/sec,  $B$  represents the bandwidth in Hz. Due to the fact that all practical channels comprise some amount of noise, it is necessary to implement communication systems such that the theoretical channel bandwidth exceeds the bandwidth of the transmitted signal by some amount and thus has some bandwidth margin. It is this excess bandwidth margin that is exploited in the physical layer auxiliary channel described in this dissertation. However, the achievable auxiliary channel bandwidth is not only limited by the transmission channel, but the transceiver architecture, which is another key factor to decide the

data rate of the auxiliary channel. In the proposed prototype transceiver, we consider implementing the auxiliary channel at relatively low bandwidth in a conventional serial wireline transceiver at a very small additional cost, which will be described in detail below.

### 3.3 Transceiver Implementation

The transceiver implementation described here is applicable to wireline baseband modulation systems. The implemented transceiver embeds the auxiliary data into the primary data using modulation that can be considered as phase modulation (PM) since the phase of the primary data stream is conditionally delayed depending upon the auxiliary data values. The transceiver simultaneously recovers both the primary and the auxiliary data at the receiving end [25].

#### 3.3.1 Transmitter Architecture

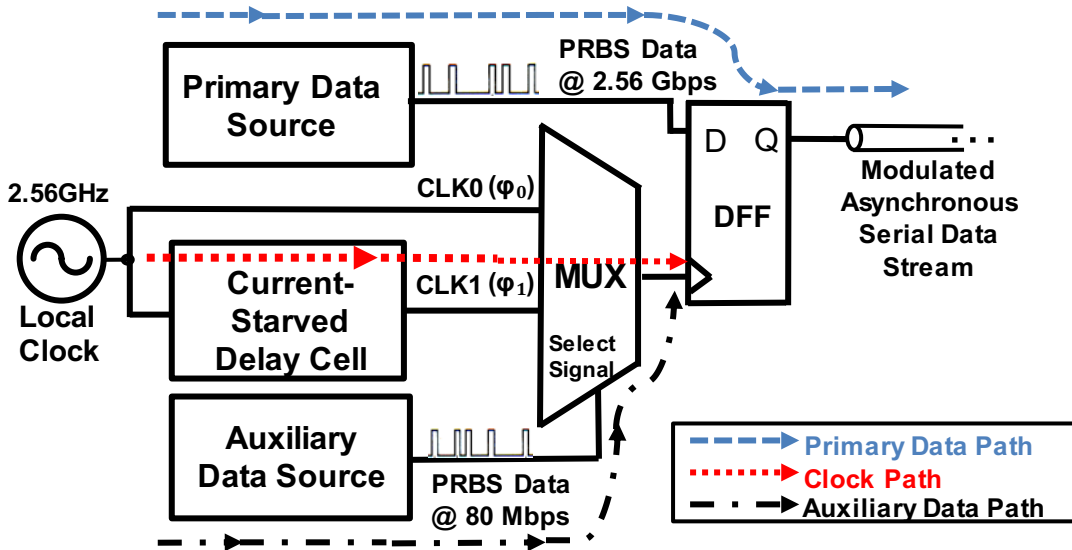
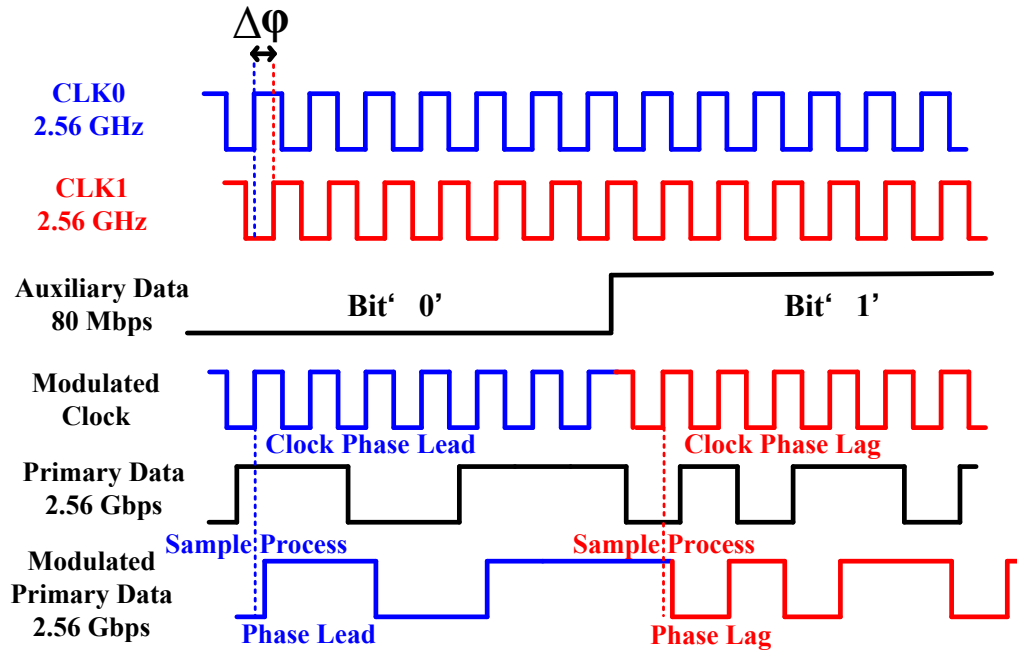


Figure 9. Block diagram of the transmitter.

A serial transmitter transmitting data at N bits per second typically employs a D flip-flop (DFF) triggered by a clock of frequency N Hz at the very end to synchronize every data bit before sending it out as serial data. The duration of every data bit is the same as that of one period of the clock. In the proposed transmitter, the auxiliary data is used to modulate the clock to the DFF. This is implemented using a 2:1 multiplexer (MUX) whose two data inputs are clock CLK0 and its delayed version CLK1. The MUX control input is driven by the auxiliary data as shown in Fig. 9. Thus, either CLK0 or CLK1 is selected for the DFF clock input by each bit of the auxiliary data signal to generate the modulated clock signal.



(a)

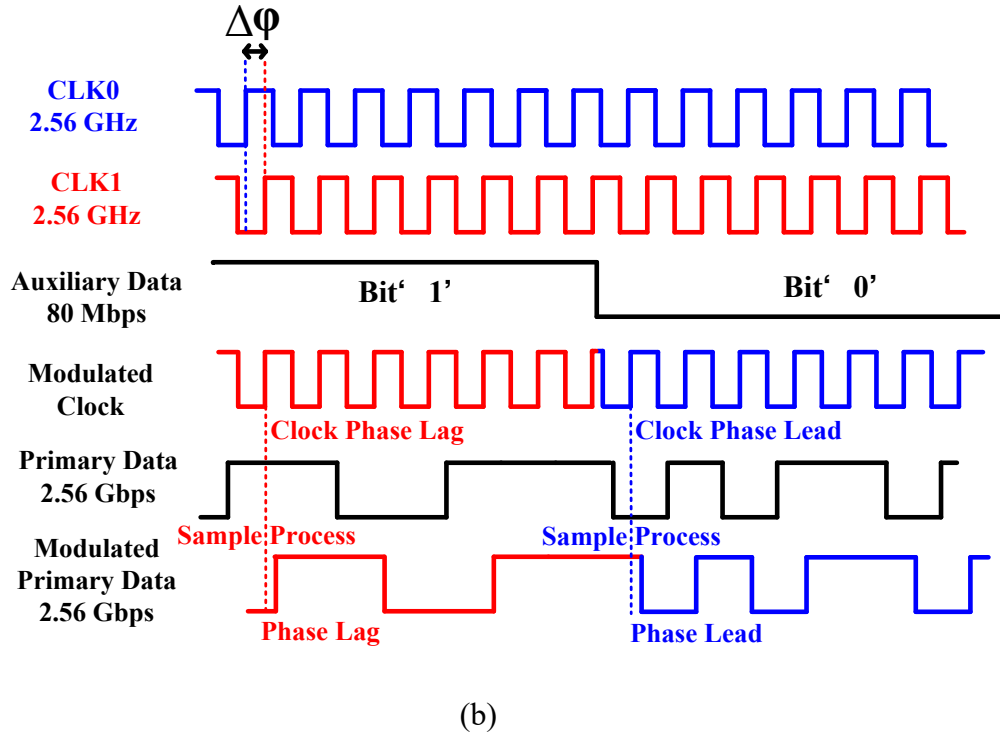


Figure 10. Timing diagrams of the transmitter modulation scheme. (a) Auxiliary data changing from bit '0' to '1'. (b) Auxiliary data changing from bit '1' to '0'.

Fig. 10 depicts the waveforms of the clocks and data in the transmitter. CLK1 is delayed from CLK0 by a phase of  $\Delta\phi$ . The instantaneous modulated clock is produced according to the auxiliary data bit. For example, when the auxiliary data bit is '0', CLK0 is used as the clock for the DFF and the modulated data stream is synchronized with the positive edge of CLK0. When the auxiliary data bit is '1', CLK1 is used for the DFF and the modulated data stream is synchronized with the positive edge of CLK1. Essentially, the '0' and '1' bits of the auxiliary data are translated to phase lead and phase lag in the modulated data stream by the MUX serving as a binary switch for CLK0 and CLK1.

Fig. 10 (a) and Fig. 10 (b) show the transmitter timing diagrams with auxiliary data changing from bit ‘0’ to ‘1’ and from ‘1’ to ‘0’, respectively. In the case from ‘0’ to ‘1’, the auxiliary data modulates the primary date from phase lead to phase lag. In the case from ‘1’ to ‘0’, the auxiliary data modulates the primary date from phase lag to phase lead.

It is necessary not to miss any clock sampling edges in the modulated clock to ensure correct data transmission. Higher data rate leads to shorter timing margin for the MUX selection process; thus the timing of the MUX block needs to be carefully verified.

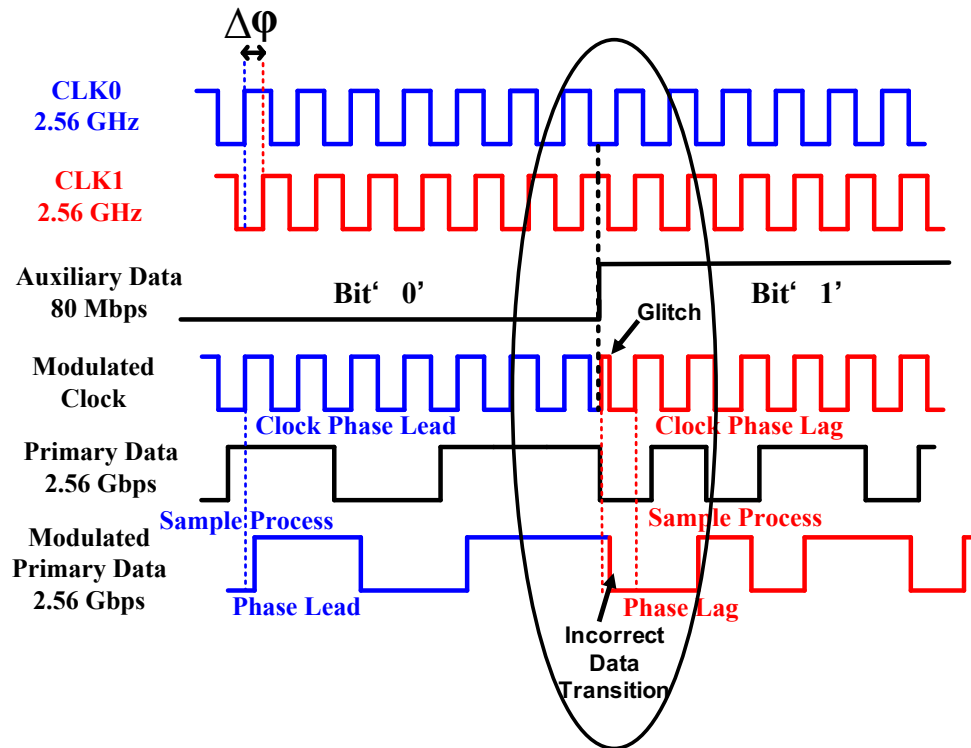


Figure 11. The case of potential glitches and incorrect data transition.



If not properly designed, glitches in the modulated clock may happen when CLK0 and CLK1 are at different voltage levels when the modulated clock makes transitions from CLK0 to CLK1 or vice versa, as shown in Fig. 11. The potential glitches are avoided by ensuring that the MUX select signal, the “auxiliary data”, always transitions when both CLK0 and CLK1 are at the same level. This is accomplished by making sure “auxiliary data” is internally synchronized to the positive edge of CLK1 (the transmitter is positive edge triggered by CLK0 and CLK1). Since CLK1 is a delayed version of CLK0 and “auxiliary data” is triggered by the positive edge of CLK1, the auxiliary data only transitions when both CLK0 and CLK1 are ‘1’, thus the potential glitches problem mentioned above would never occur.

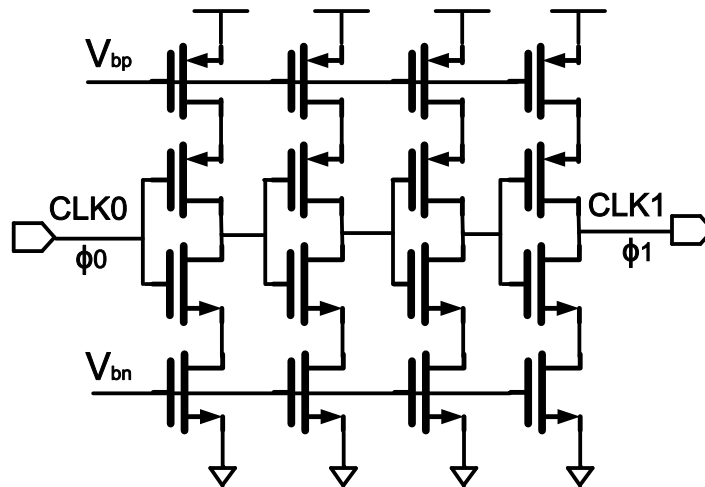


Figure 12. Current-starved delay cell.

Fig. 12 shows the current-starved delay cell that consists of four cascaded current-starved inverters. The bias voltages,  $V_{bp}$  and  $V_{bn}$ , are used to adjust the current through the top and bottom

transistors thus controlling the degree of the phase shift between CLK0 and CLK. The 10% PVT variations of the phase delay can be adjusted by the current-starved delay cell.

### 3.3.2 Receiver Architecture

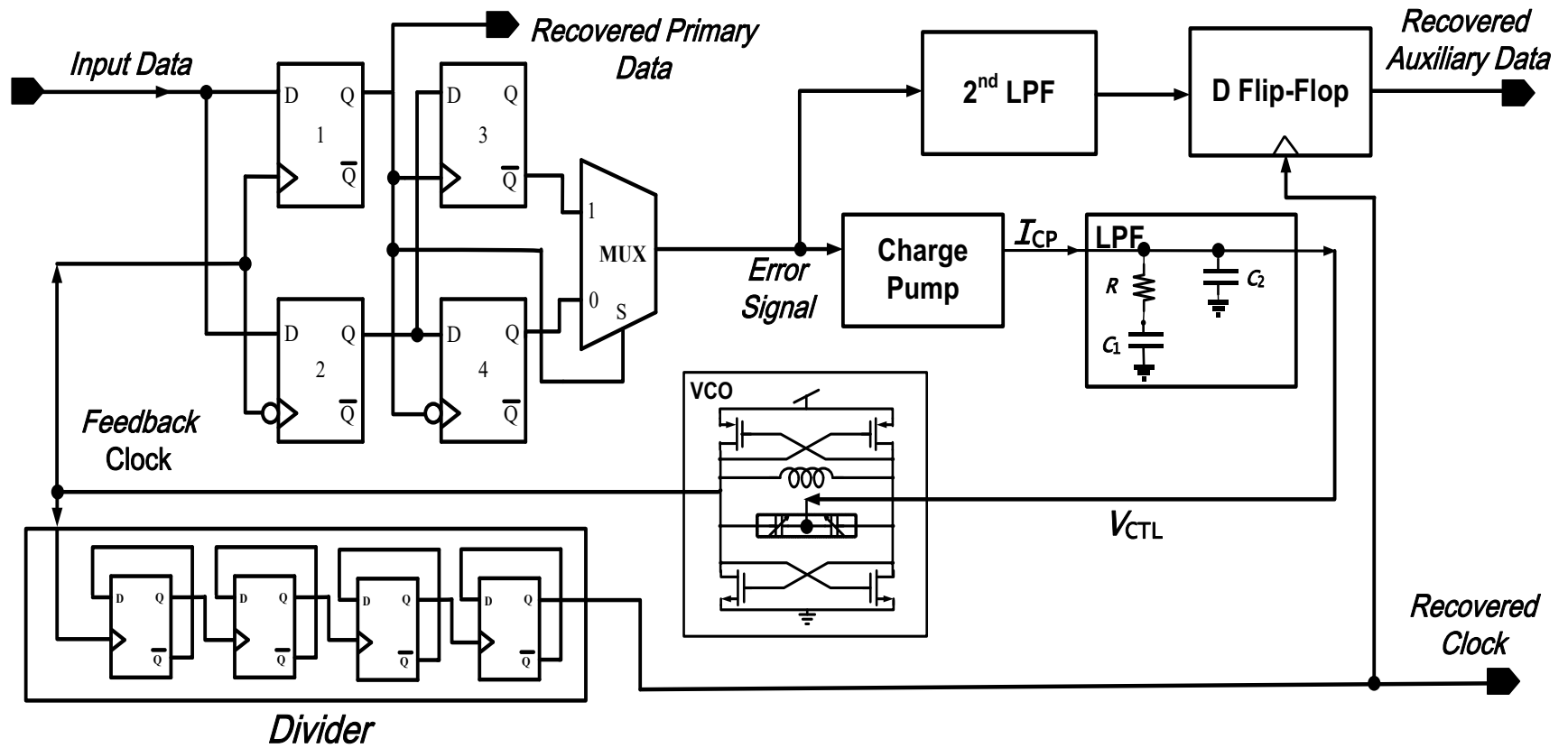


Figure 13. Block diagram of the receive

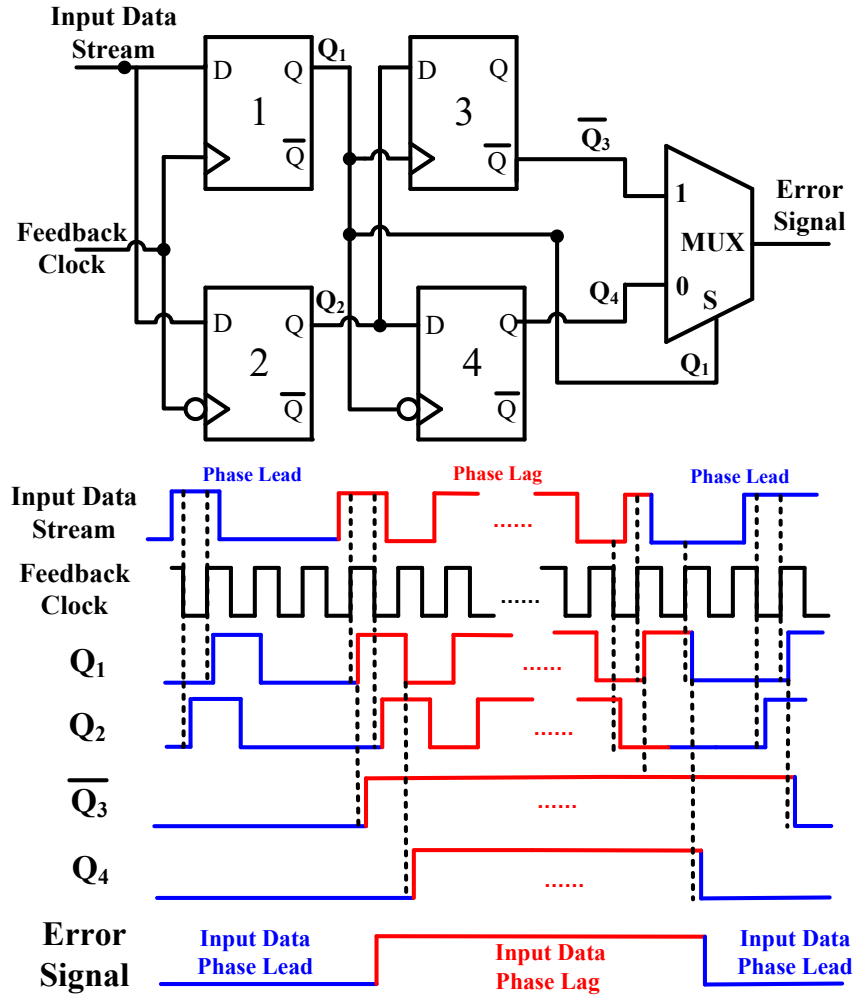


Figure 14. The architecture and timing diagram of the BBPD.

Typically, the asynchronous serial protocols require a clock and data recovery (CDR) circuit at the receiver, as the timing is embedded within the transmitted data. Fig. 13 depicts the block diagram of the receiver architecture. It contains a main CDR circuit to recover the primary data and the embedded clock, and an auxiliary data recovery path to extract and demodulate the auxiliary data from the modulated phase of the primary data. The main CDR loop is composed of

a bang-bang phase detector (BBPD), a charge pump followed by a low-pass filter (LPF), and an LC tank voltage-controlled oscillator (VCO) [26]-[28].

In particular, in the proposed receiver the BBPD is shared by both the main CDR loop and the auxiliary data recovery path. Fig. 14 shows the architecture and the timing diagram of the modified BBPD. A conventional Alexander BBPD uses three sampled points and XOR gates to produce the early or late signals that may not be continuous during the leading or lagging phase [29]. The uniqueness of the modified BBPD is that it detects the phase difference between the input data and feedback clock and produces the either '1' to indicate a phase lag or '0' to indicate a phase lead for the entire phase leg/lead period. This is an important feature to extract the auxiliary data.

The error signal produced by the BBPD is sent to the charge pump in the main CDR loop and is also used in the auxiliary data path for demodulation of the auxiliary data. In the main CDR loop, the error signal through the charge pump is used to control the VCO that adjusts the phase of the recovered clock signal so that the primary data stream can be recovered. In the auxiliary data path, as the bits '0' and '1' of the auxiliary data are embedded as phase lead and lag in the primary data stream produced by the transmitter, the BBPD error signal contains the demodulated phase information (lead/lag) that is used to recover the auxiliary data. The timing diagram in Fig. 14 indicates various signals in the BBPD for the case where the auxiliary channel is modulated either as the phase lead and phase lag in the input data. Phase lead represents bit '0' and phase lag represents '1' in the auxiliary channel. The BBPD consists of four DFFs, DFF1 through DFF4, followed by a MUX. The input data is connected to the 'D' input of the DFF1 and DFF2 whereas the VCO feedback clock is tied to the clock of the two DFFs as the sampling clock. It should be

noted DFF1 is positive-edge triggered whereas DFF2 is negative-edge triggered as shown in the figure.

During the locking process of the modulated input data stream, the BBPD error signal generated through the negative feedback loop adjusts the VCO feedback clock so that it is locked to the average phase of the input data. Once the locking status is achieved, the VCO clock is locked to the mid-point of the two phases (phase lead and phase lag) of the input data. Signal  $Q_1$ , which is always triggered at the positive edge of the VCO clock to be the recovered primary data. Signal  $Q_2$ , triggered at the negative edge of the VCO clock is either  $180^\circ$  leading or lagging with respect to  $Q_1$  and this phase relation represents the auxiliary data bits.  $Q_3$  is produced by sampling  $Q_2$  using the positive edge of  $Q_1$ , thus it produces '1' for phase leading and '0' for phase lagging.  $\overline{Q_3}$  is the opposite of  $Q_3$ . On the other hand,  $Q_4$  is produced by sampling  $Q_2$  using the negative edge of  $Q_1$ , thus it produces '0' for phase leading and '1' for phase lagging.  $\overline{Q_3}$  and  $Q_4$  are almost identical except for a certain delay between them. The final error signal is produced by selecting either  $\overline{Q_3}$  or  $Q_4$  using  $Q_1$  as the select signal to ensure bit '0' and '1' are extracted in a timely manner for the recovered auxiliary data. Although this is properly referred to as an error signal with respect to the phase of the primary channel, it is also the signal that carries the auxiliary channel bitstream and is further processed to extract the auxiliary channel bitstream.

The CDR relies upon transitions in the received data stream to produce the error signal that controls the VCO. For this reason, it is necessary that a sufficient number of transitions or pulse edges be presented in the received asynchronous data stream. In order to provide a sufficient number of signal transitions in the received primary data stream, most asynchronous serial data communications systems use encoding schemes such as pseudorandom binary sequence (PRBS)

or 8B10B encoding. These encoding schemes insert extra timing bits that guarantee at least one signal transition occurs among some number of subsequent bits. While the timing bits are redundant in terms of information content, their presence allows for transmission speeds to be increased due to the transitions occurring often enough to ensure that the receiver maintains synchronization thus enhancing overall throughput. The PRBS7 scheme is used in both the primary data stream and the auxiliary data stream in the proposed transceiver and data channels.

Some high-frequency pulses would be generated in the BBPD output, which result from the noise coupled from the input serial stream and the VCO feedback clock. The BBPD compares each transition edge of the input serial stream with that of the VCO feedback clock to detect the phase lead or lag and produces the error signal. When the instant random noise of the input serial stream and the VCO feedback clock is larger than the modulated phase difference  $\Delta\phi$ , the modulated information of phase lead/lag would be buried in the instant random jitter. At such moment, phase error detected by the BBPD is not the modulated phase information, thus some high-frequency pulses would appear in the error signal. To fully recover the auxiliary data from the error signal, a second-order low-pass filter (2nd LPF) is employed in the auxiliary data recovery path to first filter out the high-frequency pulses in the BBPD output. By choosing an appropriate bandwidth (which will be explained in the next sub-section), the 2nd LPF can filter out the high-frequency pulses of the error signal, while maintaining the recovered auxiliary data.

The auxiliary data rate is lower than the primary data rate in the proposed transceiver, but they are synchronized with each other. The clock for sampling the auxiliary data is thus a divided-down version of the faster clock in the transmitter/receiver (fast clock being the VCO output for the receiver CDR).

A standard receiver that is not equipped to demodulate the auxiliary channel will function normally and recover the primary bit stream without noticeable error due to the presence of the auxiliary channel. Because the modulated phase difference at the transmitter does not exceed allowable tolerances with respect to the requirement of CDR at the standard receiver. In this case, the phase difference modulated by the auxiliary data appears to be phase noise/jitter on the primary data channel.

### 3.4 System Design Parameters

In order to demodulate and recover the auxiliary data and to minimize the jitter in the recovered primary data, several parameters of the transceiver need to be chosen appropriately including the modulated phase difference at the transmitter, the parameters of the CDR loop, the data rate of the auxiliary channel, and the bandwidth of the 2nd LPF.

To ensure the main CDR loop operates reliably, the total jitter from the transmitter and receiver must be less than one Unit Interval (UI), otherwise the feedback clock of the CDR would not sample the input data stream correctly, thus leading to the degradation of the BER performance of the recovered primary data. To fully recover and demodulate the auxiliary data, the total jitter from the transmitter, the receiver, and the serial channel together needs to be less than the bounded deterministic jitter caused by the modulated phase difference  $\Delta\phi$  at the transmitter, otherwise the embedded auxiliary data information would be buried in the system's jitter. In our design, a jitter budget of  $<0.3$  UI is allocated to the total jitter of the transceiver and the serial channel, and a jitter budget of  $0.3-0.5$  UI is allocated to the phase difference  $\Delta\phi$ .  $\Delta\phi$  is set to be about  $135^\circ$  in our design, which can be translated into about  $0.38$  UI. A  $\pm 10\%$  PVT (process, voltage and



temperature) variation of the phase difference  $\Delta\phi$  (0.34 UI – 0.42 UI) is still within the jitter budget.

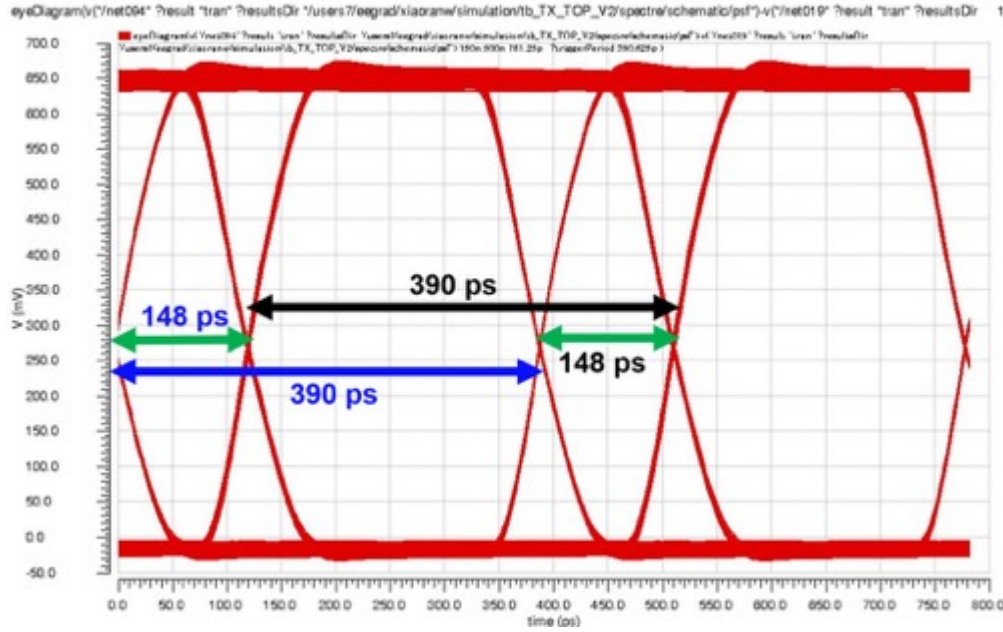


Figure 15. Simulated eye diagram of the modulated primary data.

Fig. 15 shows the simulated eye diagram of the modulated primary data with two phases. Unlike a typical transceiver eye diagram where the eye width is one UI, the proposed transmitter has the output data phase modulated by the auxiliary data and thus its eye diagram contains phase difference between bits with phase lead and bits with phase lag. The blue arrowed line marks the transmitted eye with leading phase and the black one marks the transmitted eye with lagging phase, both of which are as wide as 390ps. Note there is a 0.38 UI (148ps) phase difference, shown as indicated by the green arrows between these two eyes. This phase difference corresponds to the phase difference  $\Delta\phi$  selected for the transmitter.

The data rate of the proposed transceiver has the potential to be increased. Then the new jitter budget and modulated phase difference numbers are required to scale with the increase of the data rate. For example, if we plan to design the data rate to be ten times faster, it is necessary to scale the jitter budget and modulated phase difference to one tenth of the number we used above. It is still achievable but needs to be more carefully designed by using low-jitter techniques and some precise phase control methods. Transmitter and receiver equalizers would also be needed if the channel loss is high at higher data rate.

In a non-equipped standard receiver, the auxiliary data appears as bounded deterministic jitter of the primary data. The jitter caused by the auxiliary data still falls within the jitter budget of the transceiver, and having this much jitter would not adversely affect the functionality of the primary data recovery. Thus, the proposed auxiliary data is backward compatible with the non-equipped standard receiver.

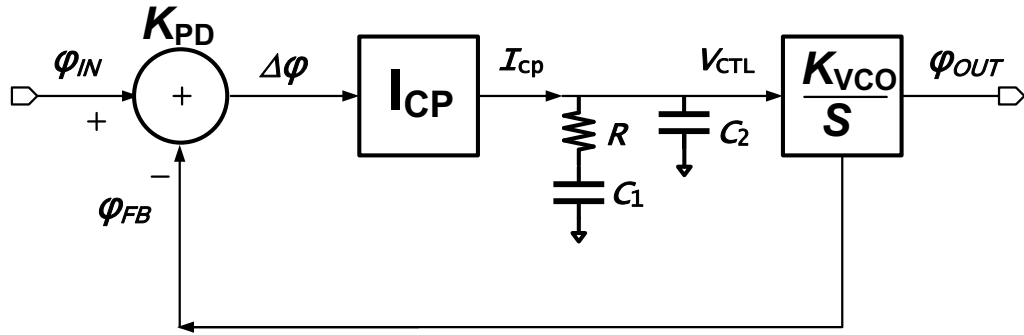


Figure 16. CDR loop linear model.

In order to select other CDR parameters, the main CDR with a negative feedback loop is modelled with a linear model as shown in Fig. 16. The small-signal model of the CDR can be derived from this block diagram.

The LPF in the CDR loop contributes to a pole at  $\omega_P$ , and a zero at  $\omega_Z$ . The magnitude frequency response of the open-loop transfer function, denoted as  $A(s)$ , is shown in (2.1).

$$A(s) = \frac{\varphi_{OUT}}{\varphi_{IN}}$$

$$= K_{PD} \cdot I_{CP} \cdot \frac{1 + \frac{s}{\omega_Z}}{s \cdot C_{total} \cdot \left(1 + \frac{s}{\omega_P}\right)} \cdot \frac{K_{VCO}}{s} \quad (2.1)$$

where  $C_{total} = C_1 + C_2$ ,  $\omega_Z = \frac{1}{R \cdot C_1}$ ,  $\omega_P = \frac{1}{R \cdot C_2}$ ,  $K_{PD}$  is the gain of BBPD,  $I_{CP}$  is the change pump current, and  $K_{VCO}$  is the gain of the VCO.

Typically,  $C_1 \gg C_2$ . (2.1) can be simplified to (2.2).

$$A(s) = K_{PD} \cdot I_{CP} \cdot \frac{R \cdot C_1}{C_{total}} \cdot \frac{K_{VCO}}{2\pi f_C} \quad (2.2)$$

where  $f_C$  is the unit gain bandwidth  $BW$  of the CDR open loop transfer function, which can be derived as shown in (2.3).

$$BW = f_C \approx \frac{K_{PD} \cdot I_{CP} \cdot R \cdot K_{VCO}}{2\pi} \quad (2.3)$$

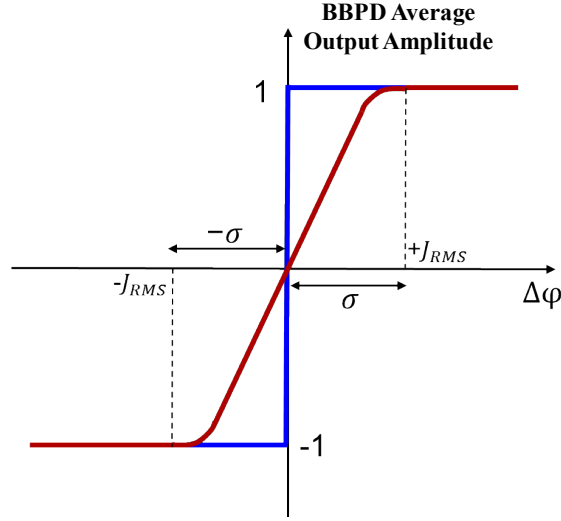


Figure. 17. BBPD output varies with the input data phase difference. Blue curve represents the ideal gain of the BBPD whereas the red curve presents the realistic gain of the BBPD.

The binary value ‘0’ and ‘1’ of the phase error signal produced by BBPD provides sign information to indicate the phase of the data is either lagging or leading the VCO feedback clock. The BBPD is followed by the charge pump to turn on either the charge/discharging current. The value ‘0’ of the BBPD output turns on the charging current of the charge pump, whereas value ‘1’ of the BBPD output turns on the discharging current. Thus ‘-1’ and ‘1’ are used to represent the BBPD output average amplitude. As shown in Fig. 17,  $K_{PD}$  is the slope of the waveform when the BBPD average output amplitude changes from “-1” to “1”. The waveform in blue color shows the ideal PD has an infinite large gain in a noiseless environment. However, in real situation, the slope of the BBPD output is influenced by the random jitter of the input data as shown in the red waveform in Fig. 17. The root-mean-square (RMS) value of the random jitter  $J_{RMS}$  in the input data stream is denoted as  $\sigma$ . The practical  $K_{PD}$  can be calculated by the slope of the red waveform times data transition density coefficient  $\alpha$ , which can be written as

$$K_{PD} = \alpha \left( \frac{2}{2\sigma} \right) \quad (2.4)$$

where  $\alpha$  is the transition density of a specific type of data input and its value depends on the average number of transition edges present during a unit time. Typically,  $\alpha = 0.5$  for the PRBS type input [30]. As mentioned above, a jitter budget of  $<0.3$  UI is allocated to the total jitter of the transceiver and the serial channel. Therefore, the maximum peak random jitter  $J_P$  cannot exceed 0.15UI, and  $\sigma$  can be calculated as

$$\sigma \approx \frac{J_P}{Q} = \frac{0.15UI}{14} = \frac{3}{280}UI \quad (2.5)$$

where  $Q$  is the scaling factor when converting between peak jitter and RMS jitter with specified  $BER$ . When  $BER = 10^{-12}$ ,  $Q \approx 14$ . Thus the value of  $K_{PD}$  can be written as

$$K_{PD} \approx \frac{1}{2} \left( \frac{1 - (-1)}{2\sigma} \right) = \frac{1}{2\sigma} = \frac{70}{3\pi} \quad (2.6)$$

The bandwidth of the CDR is set at 1.54 MHz (2.56 GHz/1667),  $K_{VCO}$  is set to  $100 \cdot 2\pi \text{ rad} \cdot \text{MHz/V}$  and  $I_{CP}$  is designed to  $1 \mu A$ . With these parameters, resistor  $R$  is derived from (2.3) as

$$R = \frac{BW \cdot 2\pi}{K_{pd} \cdot I_{cp} \cdot K_{vco}} \approx 2.1 \text{ K } \Omega \quad (2.7)$$

The value of capacitor  $C_1$  is set to 148 pF, and  $C_2$  is set to 10 pF, leading to a zero located at 512 KHz, and a pole located at 7.6 MHz as calculated below in (2.8) and (2.9).

$$f_z = \frac{1}{2\pi \cdot R \cdot C_1} = 512 \text{ KHz} \quad (2.8)$$

$$f_P = \frac{1}{2\pi \cdot R \cdot C_2} = 7.6 \text{ MHz} \quad (2.9)$$

### 3.4.1 The Auxiliary Data Rate

Shannon-Hartley Theorem reveals the theoretical bandwidth margin within the total channel capacity that can be exploited for an auxiliary channel. In the serial link design, the transceiver architecture sets the actual limitations for the auxiliary data rate, and the performance of the recovered primary and auxiliary data is related to the auxiliary data rate.

The main CDR close-loop transfer function  $H(s)$ , which is also the jitter transfer function (JTF), is expressed in (2.10), and it has a low-pass frequency response.

$$H(s) = \frac{A(s)}{1 + A(s)} = \frac{K_{PD}I_{CP}K_{VCO}R \left( s + \frac{1}{RC_1} \right)}{s^2 + (K_{PD}I_{CP}K_{VCO}R)s + \frac{K_{PD}I_{CP}K_{VCO}}{C_1}} \quad (2.10)$$

The observed jitter transfer function (OJTF) shown in (2.11) has a high-pass frequency response, so the jitter above the cut-off frequency can be observed at the output of phase detector.

$$\begin{aligned} \text{OJTF} = 1 - H(s) &= \frac{1}{1 + A(s)} = \\ &= \frac{s^2}{s^2 + (K_{PD}I_{CP}K_{VCO}R)s + \frac{K_{PD}I_{CP}K_{VCO}}{C_1}} \end{aligned} \quad (2.11)$$

Due to the low-pass characteristics of the CDR loop, the VCO feedback clock can track the low-frequency (in-band) phase change of the input data stream, and accordingly, the in-band phase noise remains at the recovered primary data, while the out-of-band phase noise appears at the error

signal of the BBPD output. In order to keep the auxiliary data information intact in the error signal, the bandwidth of the main CDR loop needs to be smaller than the lowest frequency components of the auxiliary data frequency spectrum. Otherwise, the feedback clock would track the phase change of the input data stream, which can result in the loss of embedded phase information (auxiliary data).

The random data, including the encoded data such as PRBS or 8B10B, has a wide frequency spectrum, because there are some consecutive '0's and '1's in the data stream. The PRBS7 is used for the primary and auxiliary data in the proposed transceiver, and the lowest frequency component of PRBS7 is one seventh of the data rate. In our implementation reported here, the lower limit of the data rate of the auxiliary channel is set at 7 times the CDR bandwidth. The 1.54 MHz bandwidth of the CDR leads to the lower boundary of the auxiliary data rate being 21.56 Mbps.

A ratio between the primary and auxiliary data rate is also considered to set the upper boundary for the auxiliary data rate in the proposed transceiver. As the BBPD compares the edges between the input data and the VCO feedback clock at each transition edge of the input data, the frequency of the unwanted high-frequency pulses at the error signal highly depends on the primary data rate. Therefore, the frequency range of the pulses is from the lowest to the highest frequency components of the primary data rate. In order to filter out the unwanted high-frequency pulses coupled from the noise of the primary data and the VCO feedback clock, the bandwidth of the 2nd LPF needs to be lower than the lowest frequency component of the primary data rate. Therefore, the upper limit of the auxiliary data rate that needs to be within the bandwidth of the 2nd LPF is considered to be the lowest frequency component of the primary data rate. In the proposed transceiver, the primary data is coded to PRBS7, leading to the upper boundary of the auxiliary data rate being one seventh of 2.56 Gbps, which is 365.7 Mbps.

Taking these analyses into consideration, the auxiliary data rate in the proposed transceiver is set to 80 Mbps, within the auxiliary data rate range as analyzed above. The bandwidth of the 2nd LPF is set to 40 MHz to filter out the high-frequency pulses while keeping the recovered auxiliary data.

Generally, there are some trade-offs between the performance of the recovered data (BER or jitter performance) and the auxiliary data rate. Higher auxiliary data rate that are exploited from the channel bandwidth margin may degrade the BER performance of the recovered auxiliary data. Designers can set the auxiliary data rate based on their channel SNR and the BER requirements for the received data.

The auxiliary data has some limitations and affects the performance of the primary data link on jitter tolerance and the jitter of recovered data. The jitter tolerance analysis in section III.E and jitter measurement results in section IV can depict the abovementioned limitations and impacts.

### 3.4.2 The Impact of the Auxiliary Data on the Primary Data Channel – The Jitter Tolerance Analysis

To understand the impact of the auxiliary data on the performance of the primary data channel, we evaluate the jitter tolerance (JTOL) of the primary data. JTOL can be evaluated using equation (2.12) as shown below [31].

$$JTOL(s) = \frac{\text{Timing Margin}}{\left(1 - \frac{\varphi_{OUT}(s)}{\varphi_{IN}(s)}\right)} \quad (2.12)$$



JTOL is specified in the unit of peak-to-peak jitter amplitude (UIPP). Timing margin for CDR input signal is one unit interval (1 UI).  $\varphi_{OUT}(s)/\varphi_{IN}(s)$  is the CDR close-loop transfer function  $H(s)$ . Since  $H(s)$  presents a low-pass response, CDR can tolerant the low-frequency jitter amplitude up to several times the timing margin.

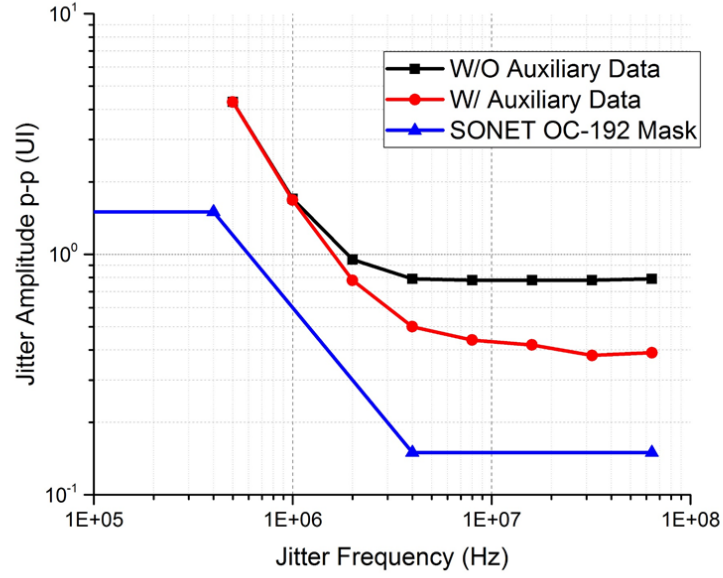


Figure. 18. Simulated jitter tolerance with SONET OC-192 mask.

The CDR jitter tolerance (JTOL) with and without the auxiliary data is simulated by varying the frequency and amplitude of a sinusoidal jitter, and measuring the maximum sinusoidal jitter amplitude, at which the receiver guarantees the targeted BER of 10<sup>-12</sup>. As can be seen from the JTOL simulation results in Fig. 18, at low jitter frequency, the JTOL performance with the auxiliary data is almost the same as the performance without the auxiliary data, whereas at high jitter frequency, the auxiliary data degrades the JTOL performance by about 0.3-0.4UI. This is

because the main CDR loop has a high-pass OJTF, and the low-frequency jitter (in band) can be tracked by CDR feedback clock. Therefore, as long as the bounded deterministic jitter introduced by the auxiliary data is within the CDR timing margin, auxiliary data does not affect the low-frequency JTOL performance. However, high-frequency sinusoidal jitter (out of band) would not be tracked by the CDR feedback clock, thus it occupies part of the CDR timing margin in the input bitstream. As mentioned above in section III. B, the auxiliary data also appears as high-frequency deterministic jitter in the CDR input bitstream and occupies some CDR timing margin. Therefore, the performance degradation of the high-frequency JTOL between the cases with and without auxiliary data is very close to the modulated phase difference  $\Delta\phi$  introduced by the auxiliary data, which is  $0.38\text{UI}$ .

The SONET OC-192 mask is also plotted in Fig. 18 for comparison. The JTOL curves of both cases (with and without the auxiliary data) exceed the SONET OC-192 mask with a margin of more than  $0.2\text{ UI}$ . This means the degraded JTOL performance is still acceptable for obtaining a BER of  $10^{-12}$  for the recovered primary data. The proposed auxiliary channel compromises some performances of the serial transceivers, like JTOL, but it still meets the communication standard and would not affect the operability.

### 3.5 Measurement Results

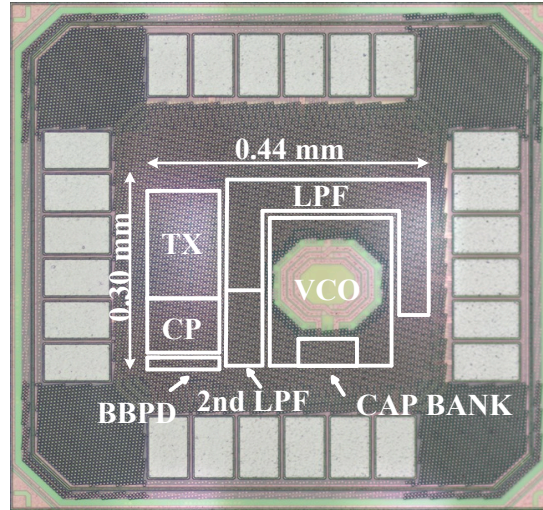


Figure. 19. Die photo of the transceiver.

The prototype IC for the proposed asynchronous serial transceiver was fabricated in a 65 nm CMOS process. The die photo is shown in Fig. 19. The core circuit occupies about  $0.13 \text{ mm}^2$ , among which  $0.03 \text{ mm}^2$  is for the transmitter, and  $0.1 \text{ mm}^2$  is for the receiver.

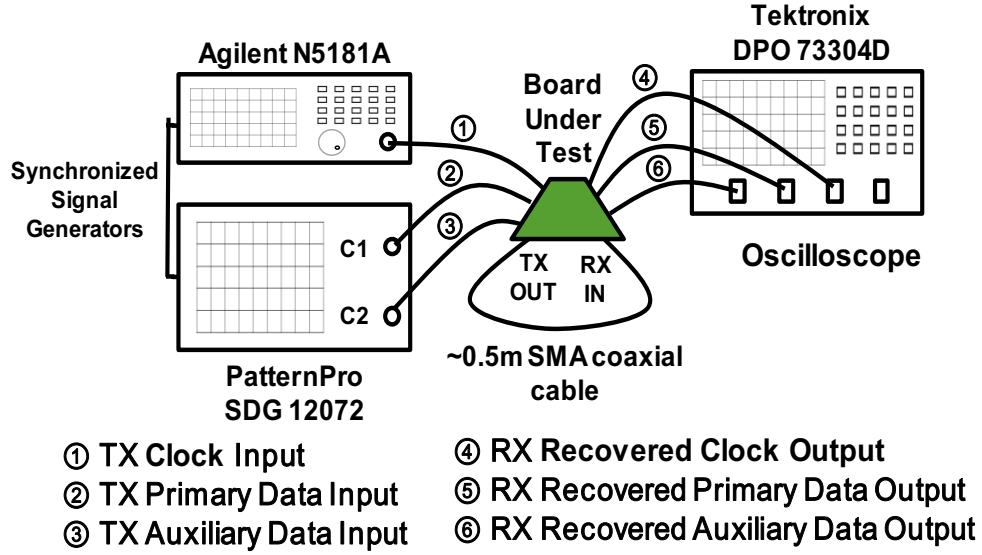


Figure. 20. The measurement environment.

Fig. 20 presents the measurement environment and connections. On the transmitter side, two synchronized signal generators (Agilent N5181A and PatternPro SDG 12072) produce three input signals, namely TX clock input, TX PRBS primary data input, TX PRBS auxiliary data input. On the receiver side, three recovered signals are connected to Tektronix 73304D oscilloscope to measure the eye diagrams and jitter performance. All external connections use short-reach (less than 0.5m) coaxial cables. The control codes were generated by an I2C master board.

The modulated data is differentially transmitted between the transmitter and the receiver by connection of a pair of differential coaxial cables. Given the 2.56 Gbps of the primary data rate and the jitter budget provided in section III. C, the  $0.38UI$  (148 ps) of the modulated phase difference with  $\pm 10\%$  variation of it (30 ps) is still within the jitter budget. Typically, the phase

delay mismatch caused by a pair of short-reach coaxial cables ( $\sim 0.5\text{m}$  in our measurement) would not exceed the margin of the phase modulation employed in the proposed design (30 ps), thus it would not adversely affect the function or performance of phase modulation and demodulation.

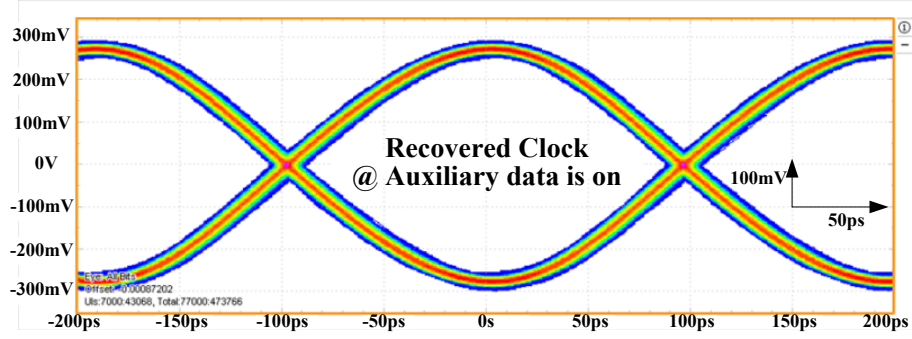


Figure. 21. The recovered clock signal.

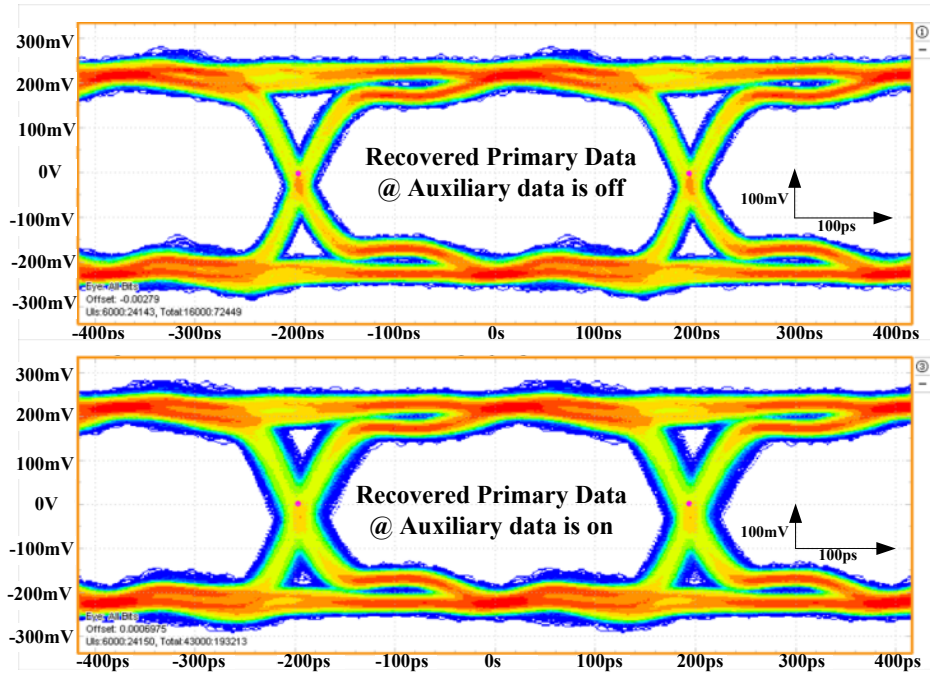


Figure. 22. The Eye diagram of the recovered primary data.

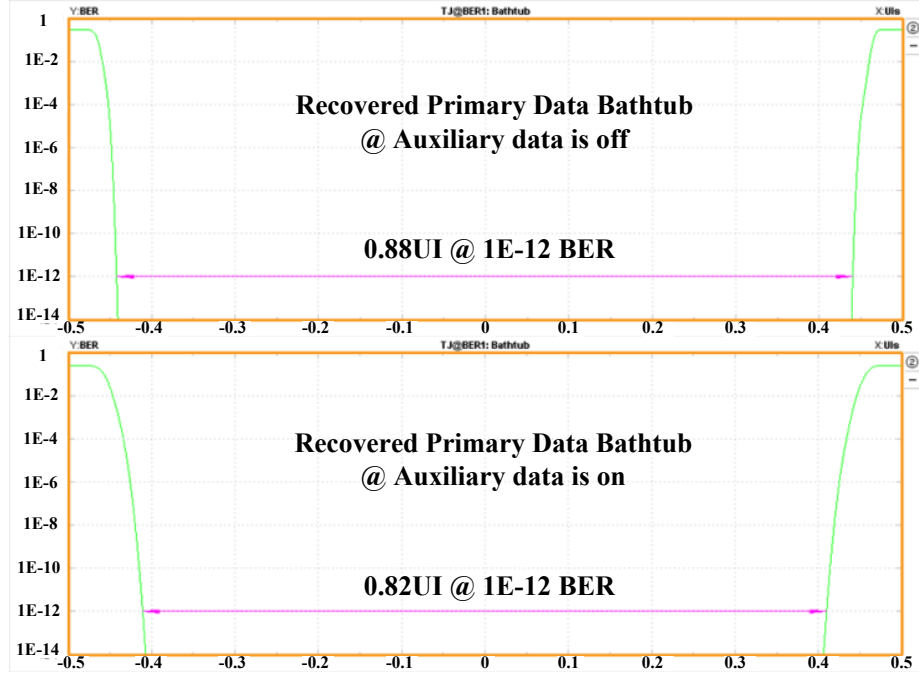


Figure. 23. The bathtub of the recovered primary data

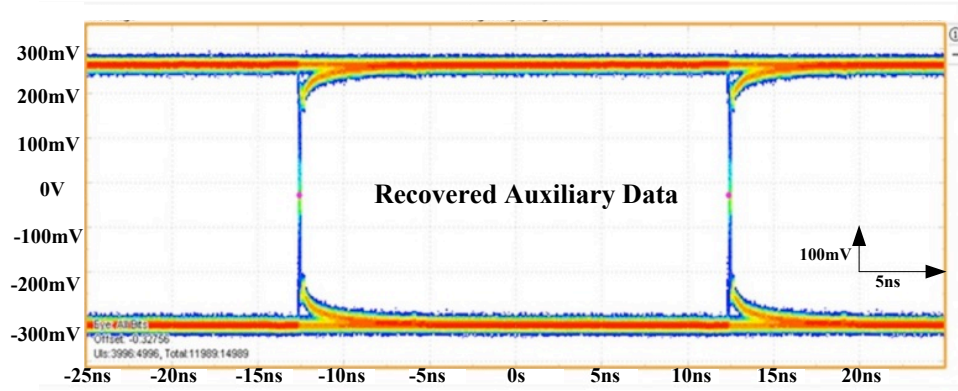


Figure. 24. Eye diagram of the recovered auxiliary data.

The transceiver operates at a supply voltage of 1.2 V. The power consumption of the core circuits is 11.8 mW at 2.56 Gbps of the primary data rate. The CDR in the receiver can cover the

primary data rate from 2.2 Gbps to 3.6 Gbps (54.7%). Fig. 21 shows the recovered 2.56 GHz clock signal with 31 ps total jitter, when the auxiliary data is in the serial link. Fig. 22 shows the eye diagram of the recovered 2.56 Gbps primary data at PRBS-7, with 48 ps total jitter when the auxiliary data is off, and with 70 ps total jitter when the auxiliary data is on. The bathtub curves of the recovered primary data shown in Fig. 23. The eye open width is 0.88 UI when the auxiliary data is off, and 0.82 UI when the auxiliary data is on. Fig. 24 shows the eye diagram of the recovered 80 Mbps auxiliary data with 103 ps total jitter. The total jitter is measured under the targeted BER of  $10^{-12}$ .

While the auxiliary data degrades the jitter performance of the recovered primary data, as can be seen from the measurement results in Fig. 22 and Fig. 23, the additional jitter (22 ps) caused by the auxiliary data is relatively small, and it would not adversely affect the functionality of the primary data recovery.

## CHAPTER 4 A HYBRID LINE DRIVER WITH VOLTAGE-MODE SST PRE-EMPHASIS AND CURRENT-MODE EQUALIZATION

### 4.1 Introduction

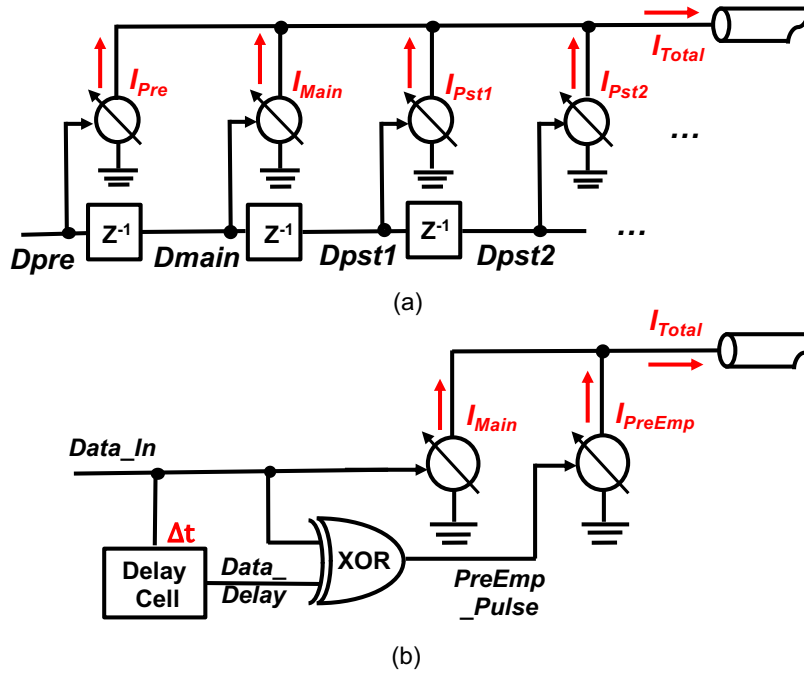


Figure. 25. Basic architecture of (a) transmitter FIR equalization and (b) pre-emphasis.

An ideal signal for a receiver completes the data transition within a symbol interval. However, signal travelling through the long cable is subject to high-frequency losses, including dielectric losses and skin effect. These losses attenuate the amplitude of high-frequency components of the transmitted signal, and the transition expands to adjacent intervals. This effect



is referred as intersymbol interference (ISI), which causes signal transmission degradation and makes it difficult to recover the signal at the receiver.

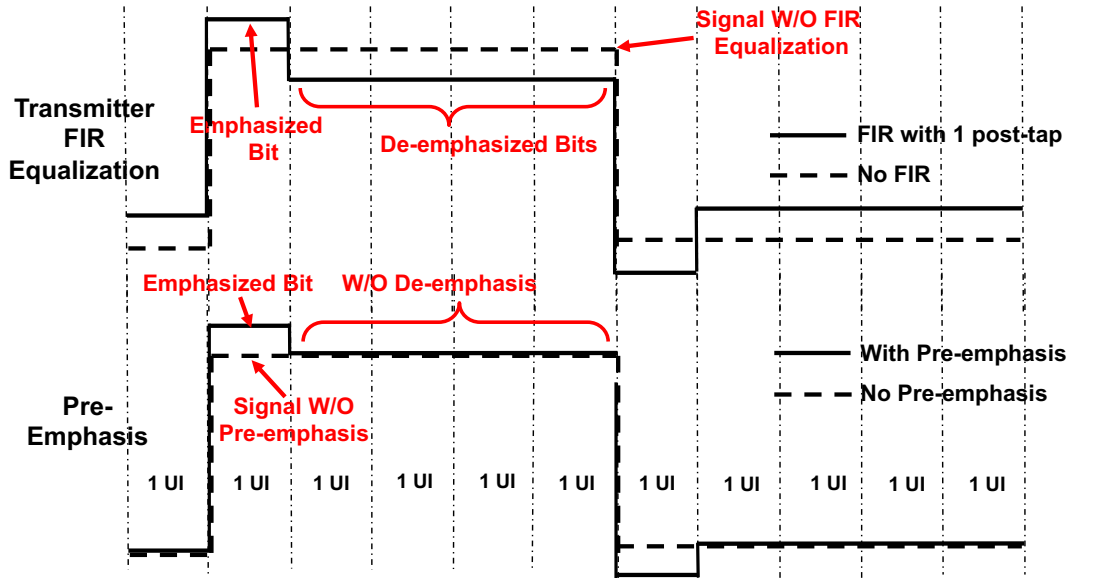


Figure. 26. Time-domain waveform of TXEQ and pre-emphasis.

In the transmitter, two ways can be used to mitigate the ISI caused by the channel loss: transmitter equalization (TXEQ) using finite-impulse response (FIR) and pre-emphasis. The FIR filter in TXEQ consists of several taps, and its function is to emphasize (boost) the high-frequency response and de-emphasize the low-frequency response to compensate the channel loss caused by ISI. It applies phase advance and delay to the signal bit and add the advanced and delayed bits to the original input signal bit with proper strength. As shown in Fig. 25 (a),  $D_{main}$  is the original input signal bit (main-cursor tap) whereas  $D_{pre}$  is the advanced bit (pre-cursor tap) and  $D_{pst1}$  and  $D_{pst2}$  are the delayed bits (post-cursor taps), with one-bit delay between the adjacent taps. The

input data streams switch on/off their corresponding current  $I_{Pre}$ ,  $I_{Main}$ ,  $I_{Pst1}$  and  $I_{Pst2}$ , which are programmable and setting the coefficients for each tap. For comparison, Fig. 25 (b) shows the basic architecture of the pre-emphasis. The *PreEmp\_Pulse* is generated by a XOR function of the *Data\_In* and its delayed signal *Data\_Delay*. *Data\_In* and *PreEmp\_Pulse* turn on/off current  $I_{Main}$  and  $I_{PreEmp}$ , respectively. Right after a transition edge of *Data\_In*, the current  $I_{PreEmp}$  is added to the nominal output current  $I_{Main}$ , which can speed up the data transition to reach the desired amplitude and decrease the rise and fall time. The duration of the  $I_{PreEmp}$  is the pulse width of the *PreEmp\_Pulse*. Pre-emphasis only boosts the high-frequency components without reducing low-frequency components, but it has limited ability to fully compensate the channel loss. The time-domain waveforms in Fig. 26 show the effects of the first post-tap in FIR equalization and the pre-emphasis. The first post-tap in the FIR equalization emphasizes the first bit period after a data transition and de-emphasizes the remaining bits. In time domain, the FIR equalization distorts the pulse shape to mitigate the spreading of the data pulses to their adjacent transmitted bits, thus eliminating ISI. However, this is at the cost of attenuating the transmitted signal. On the contrary, pre-emphasis only emphasizes the first bit period but without de-emphasizing the remaining bits, which means that pre-emphasis does not degrade the amplitude of the low-frequency components.

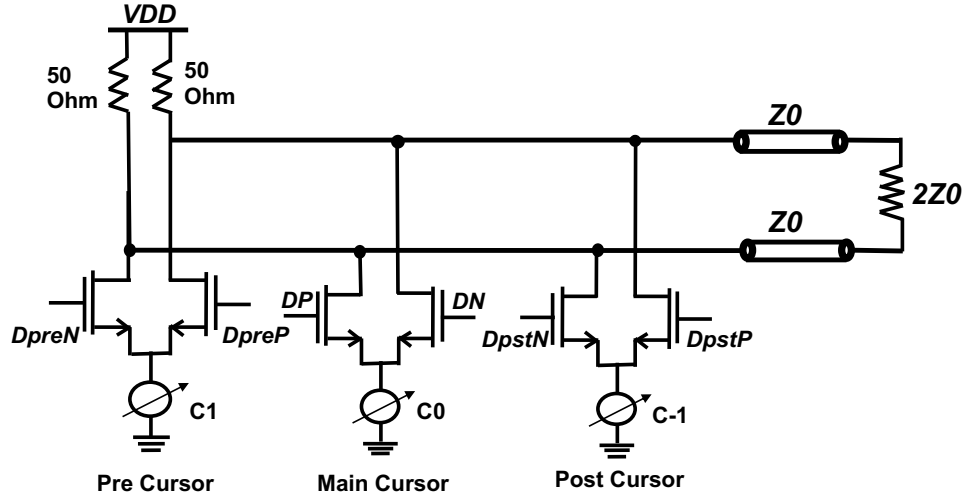


Figure. 27. Current-mode driver with FIR equalization.

A simplified schematic of a transmitter FIR equalization using CML topology is shown in Fig. 27. The pre-cursor tap, main-cursor tap and post-cursor tap are shown with the coefficient settings of C1, C0 and C-1, respectively. The tap coefficients control the current value of its tail current source to set the strength of each tap. Due to the fact that it has an inherently low susceptibility to power supply noise, the CML structure offers low jitter and low design complexity when implementing the FIR equalization at the cost of large static power consumption [34]-[34] .

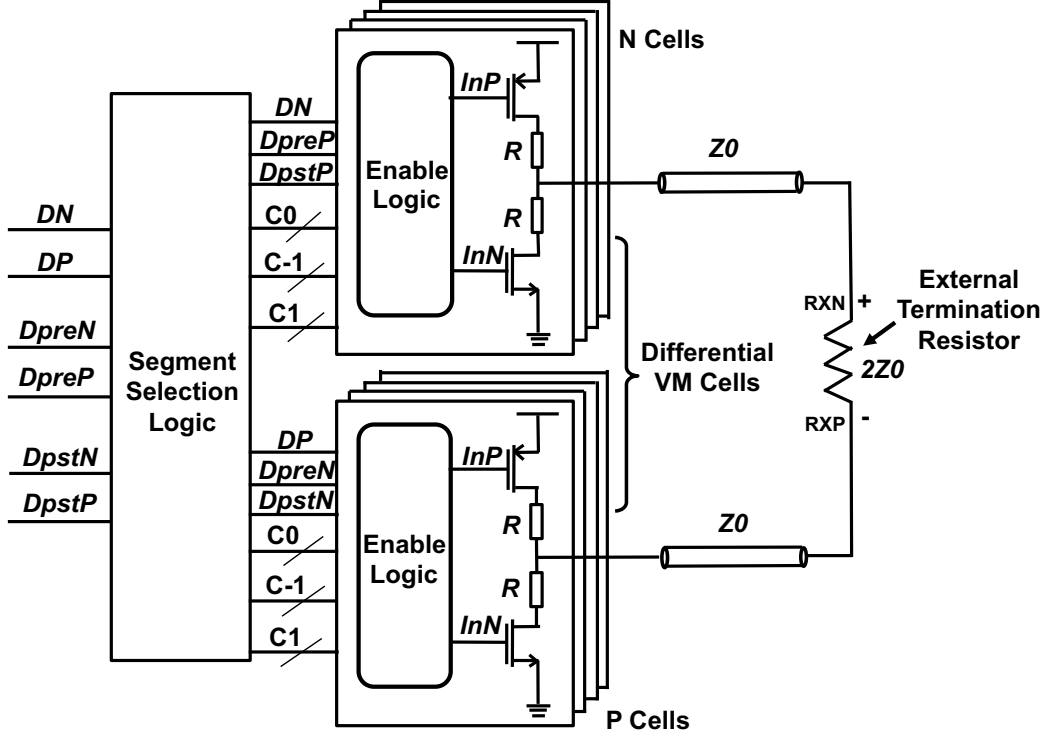


Figure. 28. Voltage-mode driver with FIR equalization.

Fig. 28 shows an SST VM driver with FIR equalization. The main-cursor signal  $DP/N$ , post-cursor signal  $DpstN/P$  and pre-cursor signal  $DpreN/P$  are the inputs for a segment selection logic. In order to implement the TXEQ in VM drivers and maintain proper impedance at the same time, the pre-driver needs an additional segment selection logic to distribute large number of segments and scale the current weight for different taps by generating coefficients  $C-1$ ,  $C0$  and  $C1$  for different group of segments [44]. This not only increases the design complexity in pre-driver, but also adds substantial capacitive load, thus degrade the power advantage of the VM drivers by increasing the dynamic power consumption especially at high data rates. Another drawback of VM

FIR equalization is that the data-dependent current would be injected into the supply rails, affecting the data-dependent jitter performance of the transmitter.

Reference [44] presented a hybrid VM main driver with CM FIR equalization. The VM main driver offers low power dissipation, and the CM FIR equalization avoids pre-driver segmentation for low design complexity. However, the amplitude of the output signal is significantly reduced by the FIR equalization, especially when very large channel loss needs to be compensated, which would be the case for the data transmission over long cables in DUNE.

This dissertation proposes a new hybrid transmitter with a line driver that combines a VM main driver with the VM pre-emphasis using SST driving cells and the CM TXEQ. The combination of the VM pre-emphasis with the CM TXEQ allows for using relatively small coefficients (strength) for the FIR equalization (thus alleviating de-emphasis of the low-frequency components) while maintaining a good eye height. Additionally, the SST-based main driver is capable of generating high signal swings with low static power consumption [45].

## 4.2 The Proposed Hybrid Line Driver

The proposed hybrid line driver is shown in Fig. 29, which consists of a differential-mode VM main driver with VM pre-emphasis and CM TXEQ as outlined in the figure. The differential input data streams, denoted as *DataN* and *DataP*, drive the VM cells, including both the main driver cells and PreEmp (pre-emphasis) cells. A 4-tap FIR equalization is implemented in the driver. The VM main driver acts as the main tap, and three CM taps (pre-cursor tap, post-cursor tap1 and post-cursor tap2) are connected in parallel with the VM main driver to implement the CM TXEQ.

The SST topology used in the VM main driver cells and VM PreEmp cells is also shown in Fig. 29. An Enable Logic generates *Up* and *Down* to switch on and off the VM driving cells. Since SST is a CMOS-oriented design, the SST cells can be turned off when the signal *Up* is set to '1', and *Down* is set to '0'. Signal *Enable* serves as the control signal for the Enable Logic in the main driver cells; both signals *Enable* and *PreEmp\_Pulse* serve as the control signals for the Enable Logic in the PreEmp cells. For example, the PreEmp cell is turned on when both its corresponding *Enable* and *PreEmp\_Pulse* are at logic '1'.

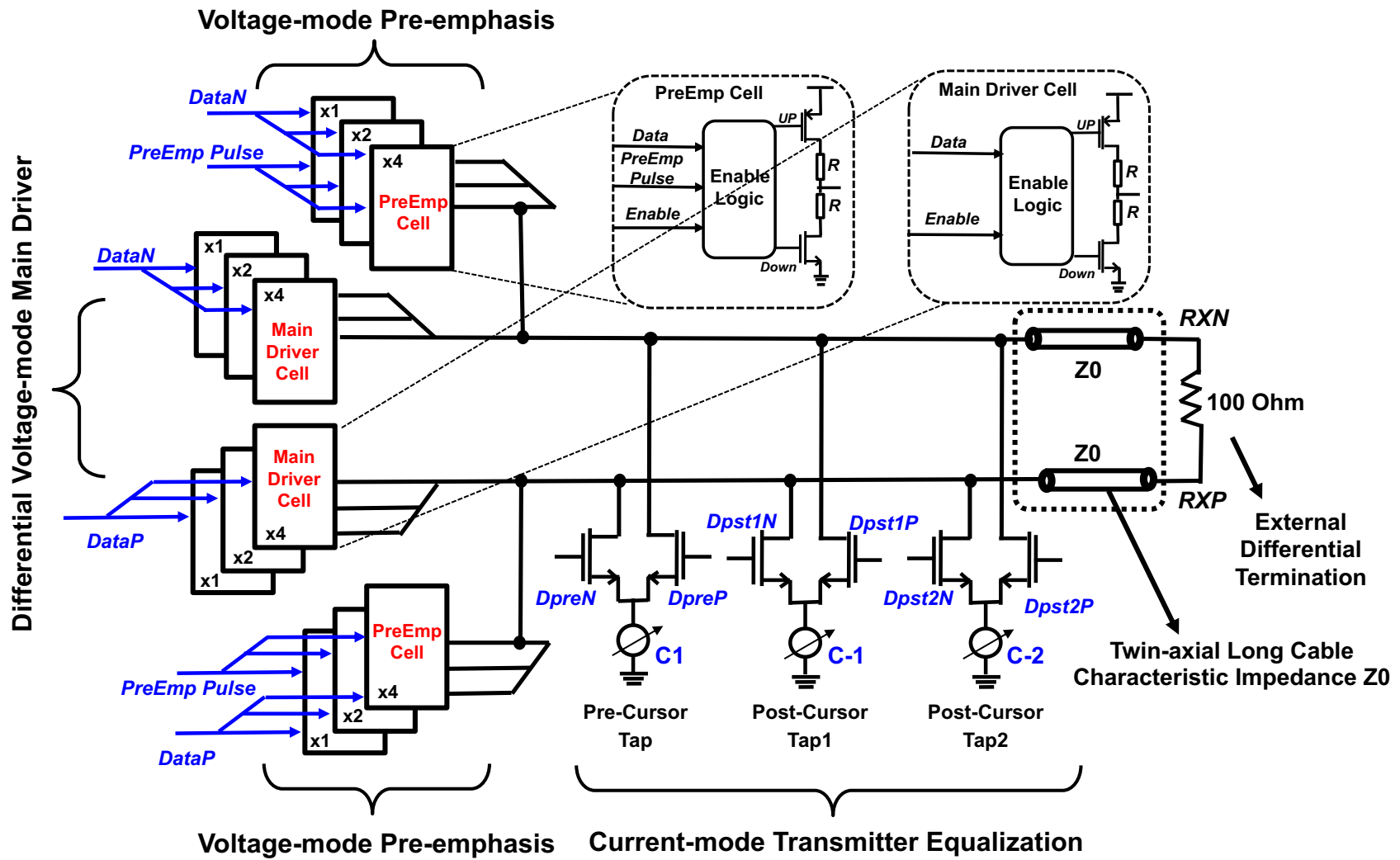


Figure. 29. The proposed hybrid line driver with voltage-mode pre-emphasis and current-mode transmitter equalization.

In the proposed hybrid line driver in Fig. 29,  $DpreN/P$  are differential inputs of the pre-cursor tap, and they lead  $DataN/P$  by one bit time.  $Dpst1N/P$  and  $Dpst2N/P$  are inputs of the post-cursor tap1 and post-cursor tap2, which lag  $DataN/P$  by one bit and two bits, respectively. The tap coefficients are implemented by tail current source with 4-bit binary control. Higher current resolution can be achieved by increasing the number of bits. However, TXEQ deals with channel loss and ISI by sacrificing the signal amplitude, which results in a trade-off between the signal amplitude and SNR/jitter performance of the transmitted signals.

In DUNE, large signal swing is required for the receiver to recover the transmitted data with a good BER performance. In order to reduce the signal attenuation caused by FIR equalization and get large eye height, the VM SST pre-emphasis is proposed in the hybrid line driver. It peaks the voltage of the first bit after the transition edges to decrease the rise and fall time of the transmitted data, which enhances the transmission bandwidth and boosts the high-frequency response. VM pre-emphasis allows the FIR equalization to use smaller tap coefficients than without pre-emphasis, thus alleviating the de-emphasis to obtain larger signal swing. The combination of CM FIR equalization and VM pre-emphasis allows the line driver to take advantage of both topologies to compensate the large channel loss while still maintaining large signal magnitude.



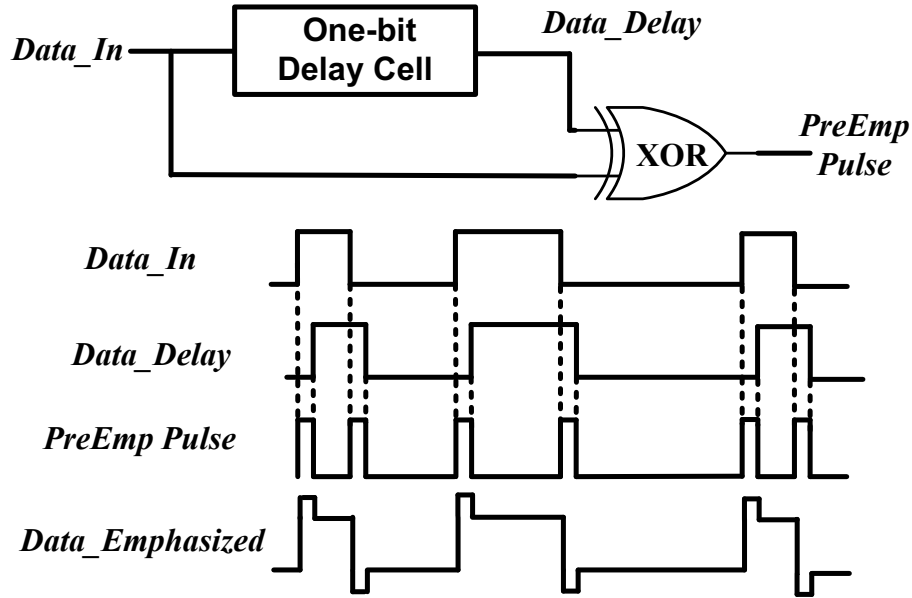


Figure. 30. “PreEmp\_Pulse” generation.

Fig. 30 shows the simplified timing diagram to depict the generation of the *PreEmp\_Pulse*. An XOR with two inputs of *Data\_In* and its delayed stream *Data\_Delay* is used to generate the *PreEmp\_Pulse*. The one-bit delay cell is used to delay the input data with one-bit period, which determines the width of the *PreEmp\_Pulse*. Logic ‘1’ of *PreEmp\_Pulse* turns on the corresponding VM PreEmp cell, and provides extra current during the pre-emphasis period, thus generating a voltage peak at the transmitted data pulses. As can be seen from the timing waveforms in Fig. 30, pre-emphasis is enabled by *PreEmp\_Pulse*, and in *Data\_Emphasized*, the voltage peak is applied to the bits after the transition edges of transmitted data.

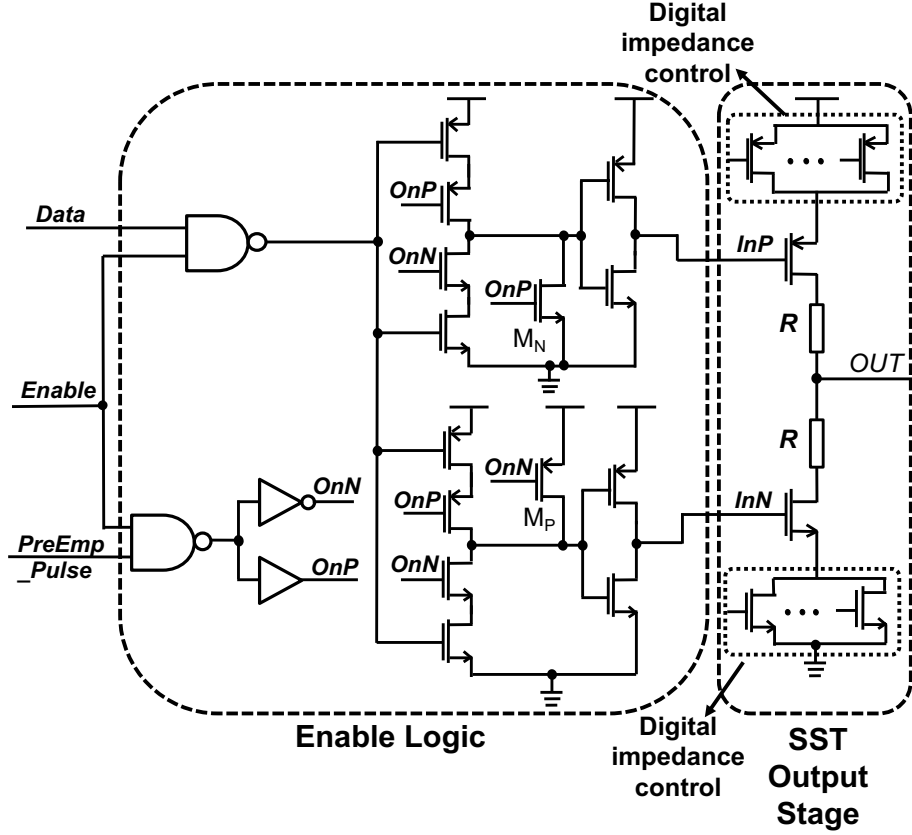


Figure. 31. The enable-logic and SST output stage in main driver cell and PreEmp cell.

The schematic of the VM main driver cell and VM PreEmp cell is shown in Fig. 31, which includes the Enable Logic and the SST output stage. The Enable Logic consists of some combination logics and two stages of inverter-type buffers for differential path. The function of the Enable Logic is to switch on and off the corresponding SST output stage controlled by *Enable* and *PreEmp\_Pulse* signals. The VM driving cells are turned on when both *Enable* and *PreEmp\_Pulse* are high. When either one of *Enable* or *PreEmp\_Pulse* is low, *OnN* and *InP* are set to low, and *OnP* and *InN* change to high, which shutting off the current flow and turning off the data switching in both Enable Logic and the SST output stage. Since the main driver cells are only

controlled by *Enable* signal, the *PreEmp\_Pulse* inputs are always set high for the main driver cells. For good jitter performance, it is important to maintain fast rise and fall time for all dynamic nodes inside the Enable Logic and the SST output stages.

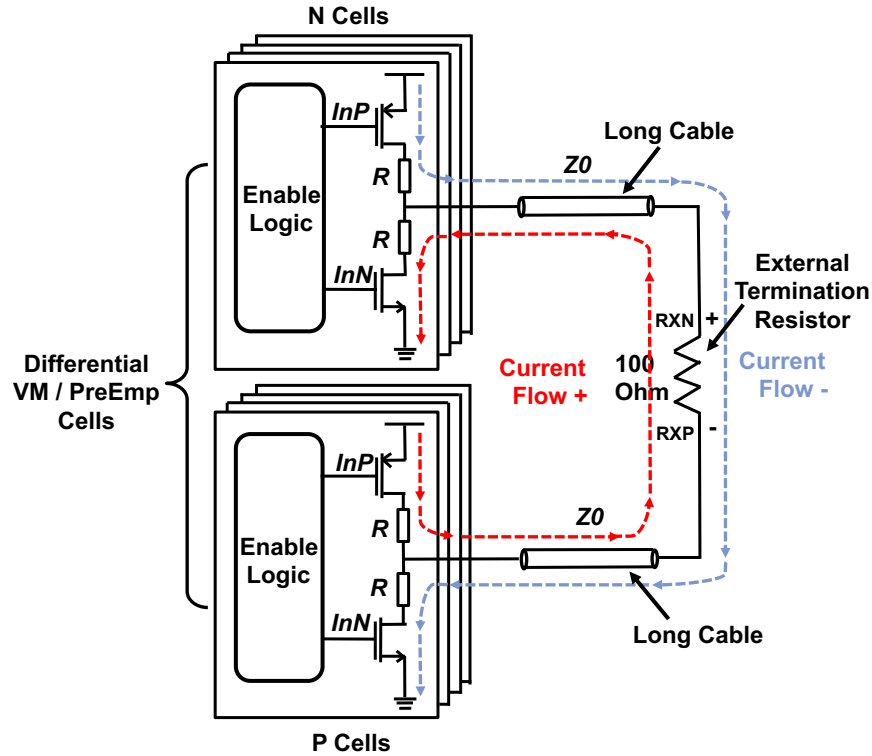


Figure. 32. The current flow in VM driver using SST output stages.

Fig. 32 depicts the current flow on differential VM driving cells, which determines the differential output signal swing of the VM drivers. The blue current arrows represent the case of input data bit is “0” for N cells and “1” for P cells, when the switching transistors of PMOS in N cells and NMOS in P cells are on. The current flows from the VDD of N cells and through two series resistors R, twinax cables and 100 Ohm external termination resistor. Similarly, the case of input data bit “1” for N cells and “0” for P cells is presented by red arrows for current flow. The

single-ended voltage swing at the receiving end is actually the voltage across the 100 Ohm external termination resistor. The output resistance of the VM driver is set to be 50 Ohm, therefore, the differential output signal swing of the VM driver is  $2 \cdot \frac{100}{2R+100} VDD$ , which is 1.1 V in the design.

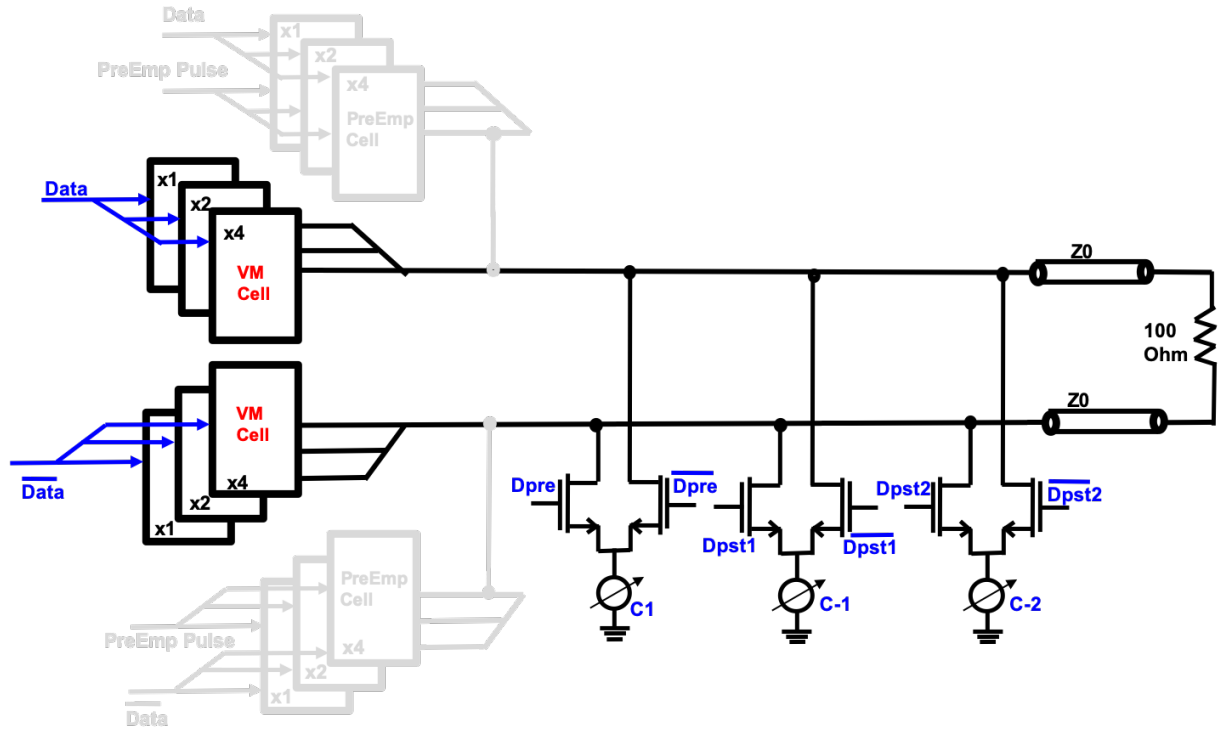


Figure. 33. Hybrid line driver without pre-emphasis.



### 4.3 DUNE Transmitter Design

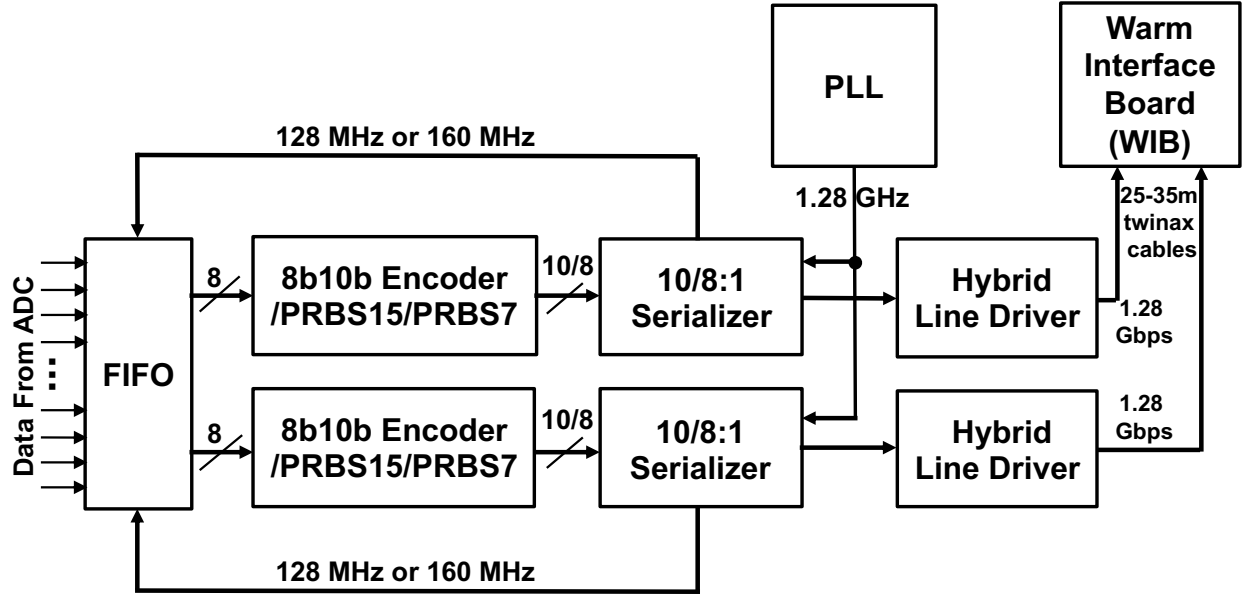


Figure. 35. The block diagram of transmitter.

Fig. 35 shows the DUNE transmitter which is a data concentrator that captures the incoming digital data streams from several ADCs and sends data from cryogenic environment (liquid argon) over long twinax cables to the warm interface. The transmitter consists of two channels of 8b10b encoder, an 8/10:1 serializer and a line driver. The phase locked loop (PLL) generates a high-speed clock to synchronize the transmitter [46]-[47]. The transmitter implements standard 8b10b encoding, and encoding can also be switched off and replaced with PRBS7/15 to facilitate debugging.

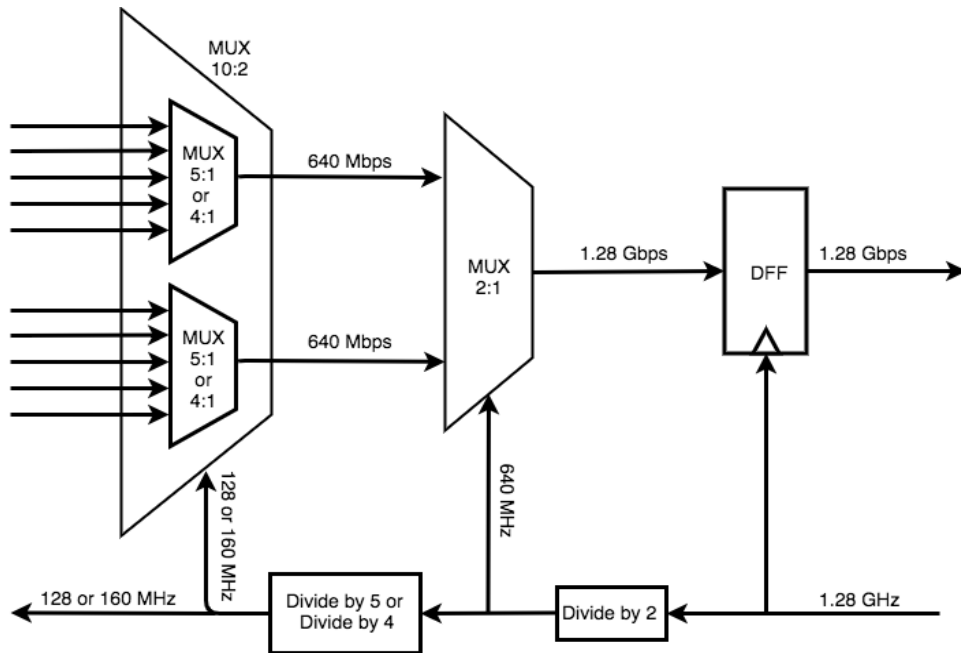


Figure. 36. The block diagram of the serializer.

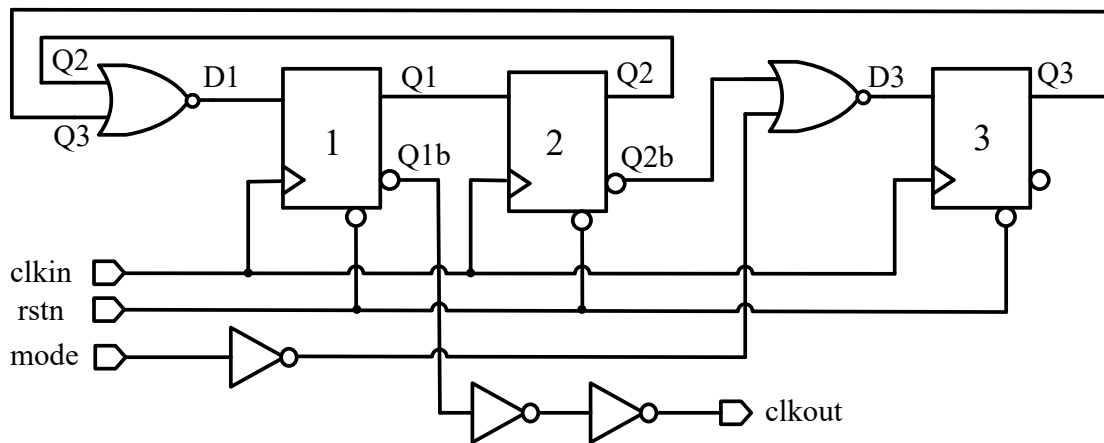


Figure. 37. The block diagram of the 4/5 divider.

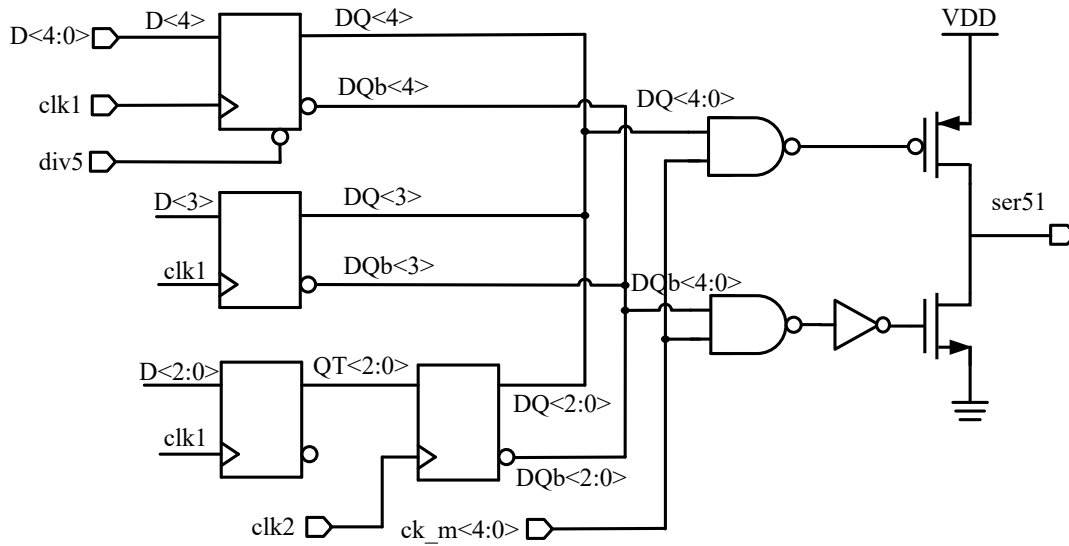


Figure. 38. The block diagram of the 4/5-to-1 MUX.

A block diagram of the serializer, the schematic of the divide-by-4/5 divider and 4/5:1 MUX are shown in Fig. 36, Fig. 37 and Fig. 38, respectively. In the normal 10-bit mode, incoming parallel data at 128 MHz is input to two 5:1 multiplexers, whose output is then input to a 2:1 multiplexer, registered by a D flip-flop, and output to a line driver. The 640 MHz and 128 MHz clocks required by the multiplexers are produced in each serializer stage. The 128 MHz clock is also to the FIFO, which provides synchronized data to the serializer. An 8-bit mode is included for debugging purposes. In 8-bit mode, parallel data is input at 160 MHz and the first multiplexer stage operates as two 4:1 multiplexers.



#### 4.4 Design Considerations for Lifetime Reliability

It is required that electronics including integrated circuits (ICs) used in DUNE have at least 20 years of the experiment lifetime. The read-out circuits including the proposed transmitter will be immersed in liquid argon at cryogenic temperature (89K). The challenge is to design the transmitter to operate at such low temperature with the required lifetime.

Some of the transistor performance such as thermal noise, speed, driving current can be improved at cryogenic temperature. However, at lower temperature, hot carriers in the CMOS devices generate larger energies with higher frequencies. Therefore, the hot carriers would degrade the transistor lifetime at cryogenic temperature, which is a primary concern for DUNE electronics design.

Our previous research carried out the lifetime study on commercial 130nm and 65nm CMOS processes [48]-[52]. It was reported that the 65 nm process is more resistant to cryogenic hot carrier degradation than the 130 nm process. The transistors in 130 nm process are required to reduce its nominal voltage from 1.5 V to 1.49 V in order to get 20 years of lifetime. While for 65 nm process, the voltage below 1.3 V meets the 20 years lifetime requirement. Given the fact that the lifetime is inversely proportional to the supply voltage of the transistors, to further push the lifetime up to 30 years, a 1.1 V supply voltage of 65 nm CMOS process is determined to be used in the DUNE cold electronics, including the proposed DUNE transmitter.

In addition, the studies also showed that lifetime increases with increasing transistor lengths, but remains unaffected by the transistor width. The width dependence is negligible for short channel devices. After balancing the trade-offs between lifetime and other transistor performance,

such as speed, 90 nm instead of 65nm minimum transistor length is chosen to be used in the DUNE cold electronics IC design.

#### 4.5 Measurement Results

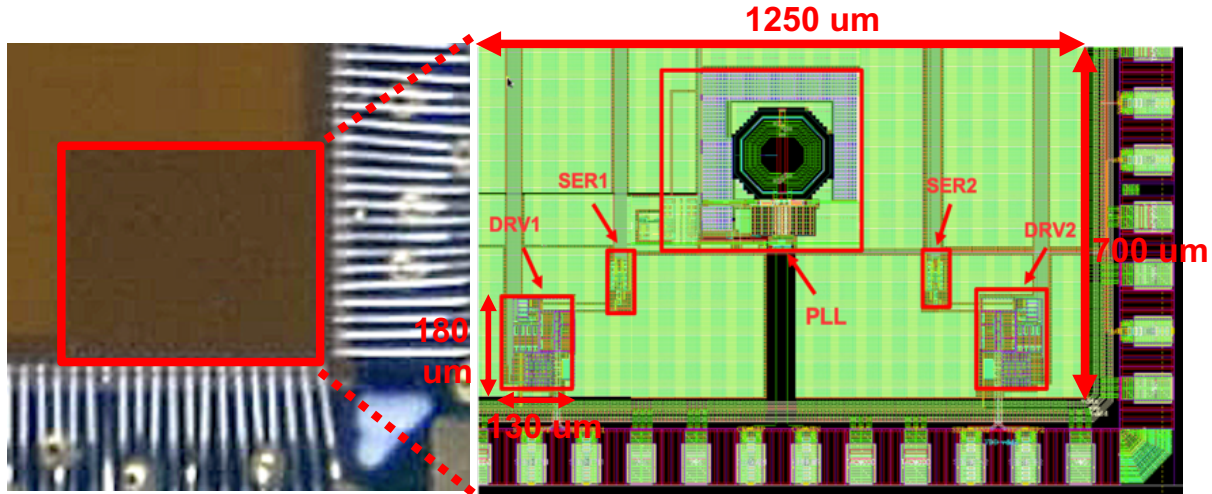
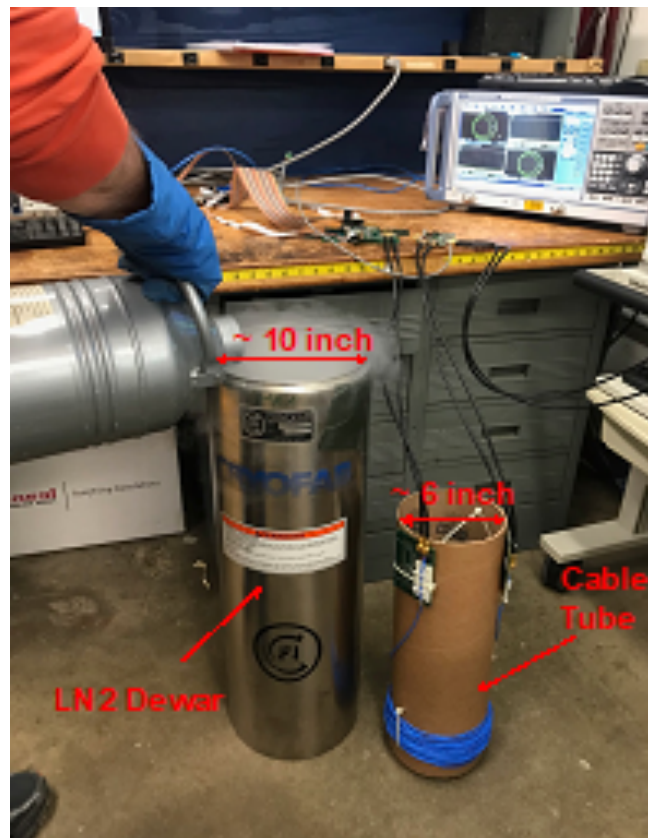


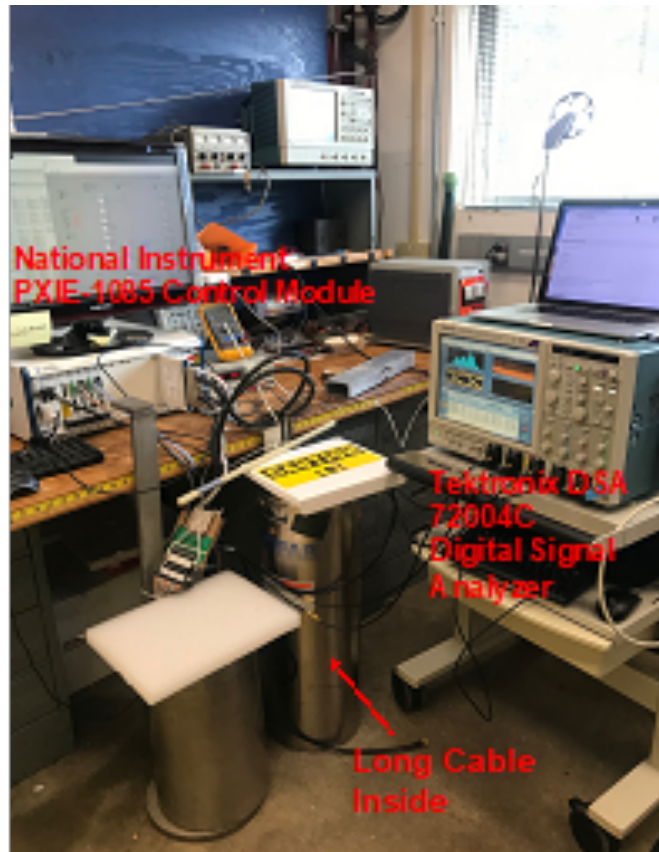
Figure. 39. Chip die photograph and corresponding layout view.

The proposed hybrid line driver is designed and fabricated in a 65 nm technology. The front-end transmitter is implemented with entire ASIC. The die photo of a part of the chip and the corresponding layout of the transmitter are shown in Fig. 39. The transmitter has two channels of data transmission which share one phase-locked-loop (PLL). Therefore, the transmitter layout contains one PLL, two channels of serializers (SER1 and SER2), and two channels of proposed line drivers (DRV1 and DRV2). The core area of the proposed hybrid line driver takes about  $180 \mu\text{m} \times 130 \mu\text{m}$ . For easy measurements, liquid nitrogen with the temperature of 77 Kelvin is used,

which is close to that of liquid argon (89K). The proposed hybrid line driver is tested by immersing the entire chip into liquid nitrogen. Its ability to drive 25 m and 35 m cables is verified for both warm cable (300K) and cold cable (77K). Measurement results indicate that three different driver configurations, namely, the proposed hybrid line driver, hybrid driver without pre-emphasis, and CML driver with only CM TXEQ are able to drive the long cables. PRBS7 and PRBS15 data patterns are generated and transmitted, both data patterns and their pattern length are received and recognized by Tektronix DSA 72004C oscilloscope for three above-mentioned driver configurations.

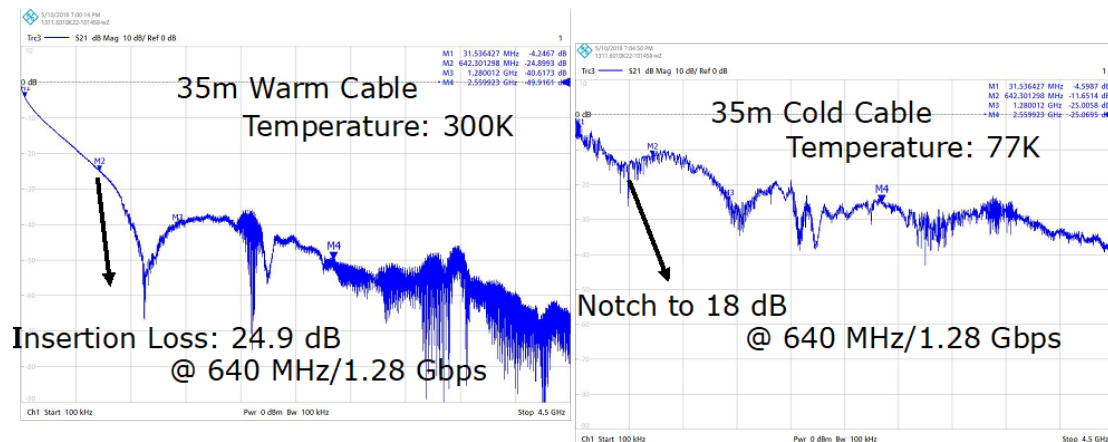


(a)



(b)

Figure. 40. Test set up and environment for line driver measurement.



(a)

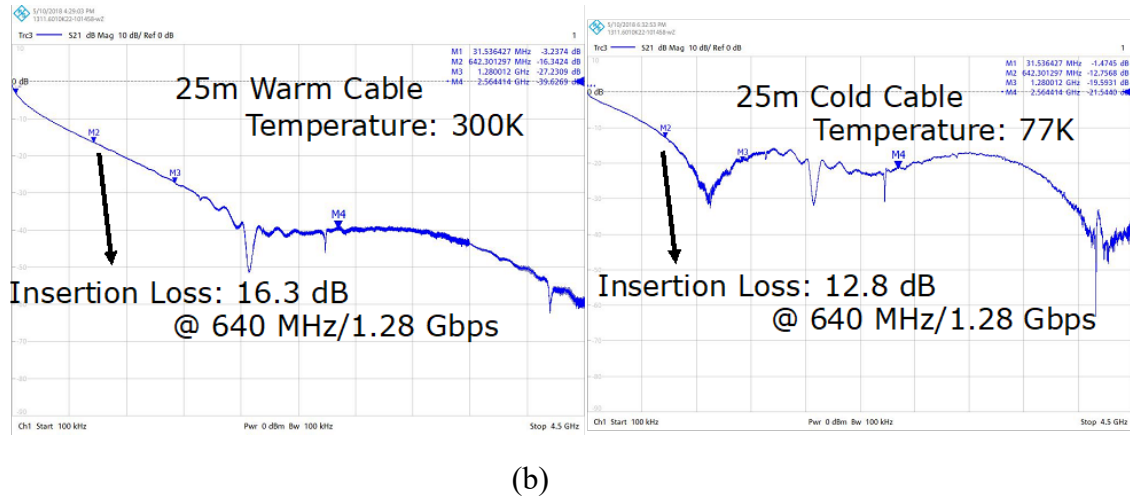


Figure. 41. Insertion loss of (a) 35m and (b) 25m twinax cable at room and cryogenic temperature.

Fig. 40 shows the test set up and environment for the line driver measurement. As can be seen from left picture, the Samtec twinax cable (blue color) is twined at the bottom of a cable tube with the diameter of about 6 inches. When testing with the cold cable, the cable tube is put into a large liquid nitrogen dewar, and the long cable is completely immersed into liquid nitrogen to be cooled down to cryogenic temperature, as shown in Fig. 40 (b). The output of the line driver is connected to one end of the long cable, and the other end of the long cable is connected to the Tektronix DSA 72004C oscilloscope without any equalizer at the receiving end. Fig. 41 (a) shows the insertion loss of a 35-meter cable is 18 dB at 77K and 24.9 dB at 300K, and Fig. 41 (b) shows the insertion loss of a 25-meter cable is 12.8 dB at 77K and 16.3 dB in 300K.

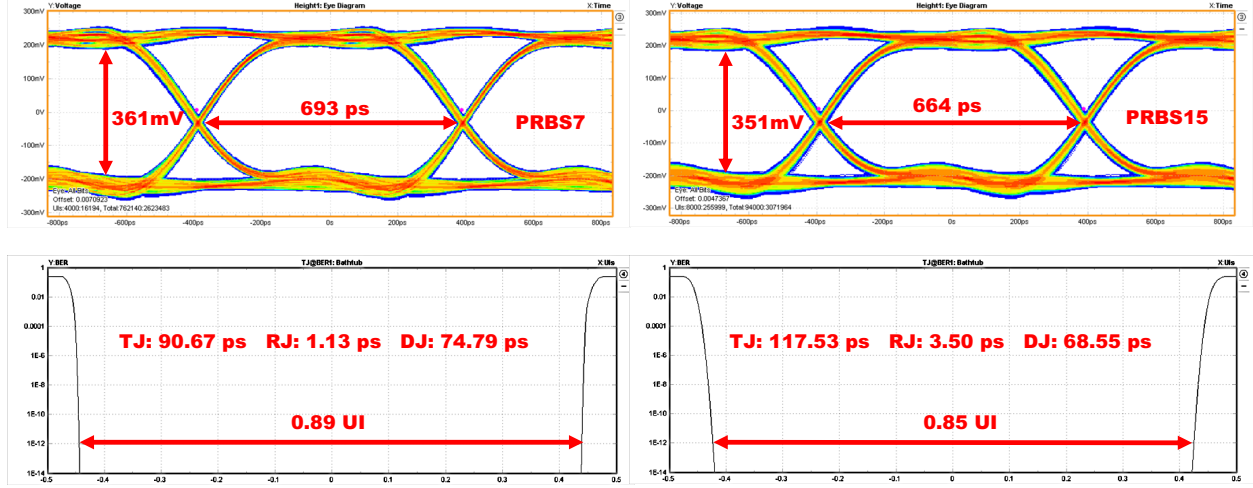


Figure. 42. The measured transmitted eye diagrams and bathtub curves of the serializer output for input data of (a) PRBS-7 and (b) PRBS-15.

Fig. 42 shows the measured transmitted eye diagrams and bathtub curves of the serializer output signals, which are obtained by programming the driver to a conventional CML driver that drives a differential pair of short coaxial cables. PRBS-7 input data pattern obtains 361 mV eye height and 693 ps eye width (0.89UI). PRBS-15 input data pattern obtains 351 mV eye height and 664 ps eye width (0.85UI). The total jitter (TJ), random jitter (RJ) and deterministic jitter (DJ) for PRBS-7 are 90.7 ps, 1.1 ps and 74.8 ps, respectively. For PRBS-15 data pattern, they are 117.5 ps, 3.5 ps and 68.6 ps, respectively.

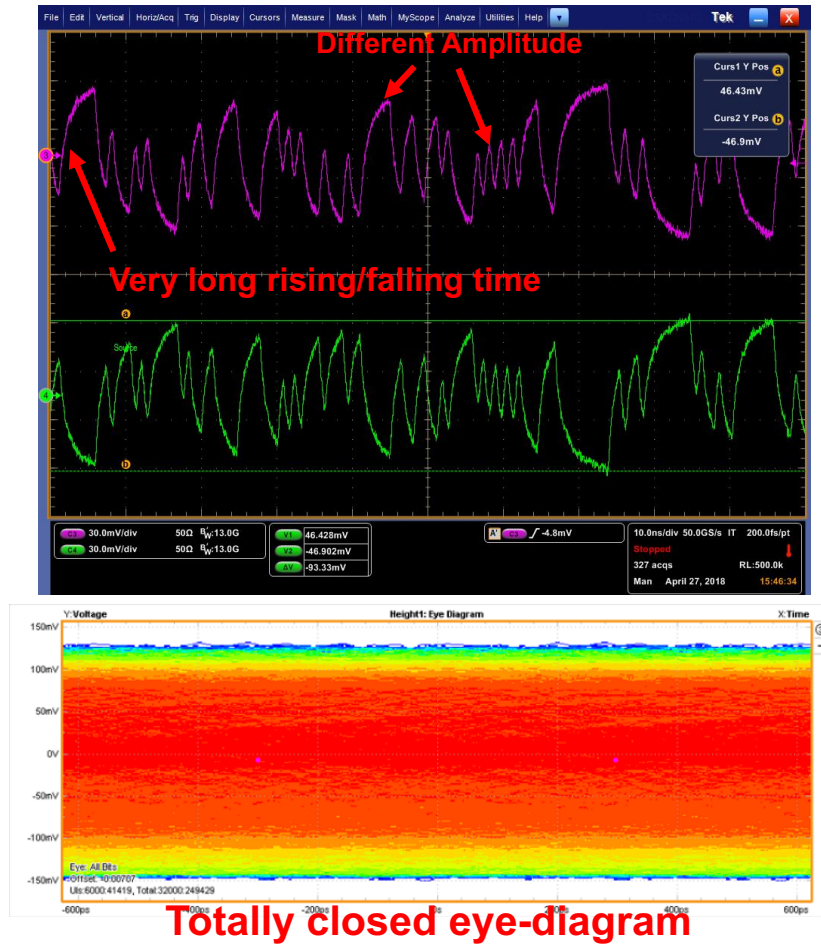
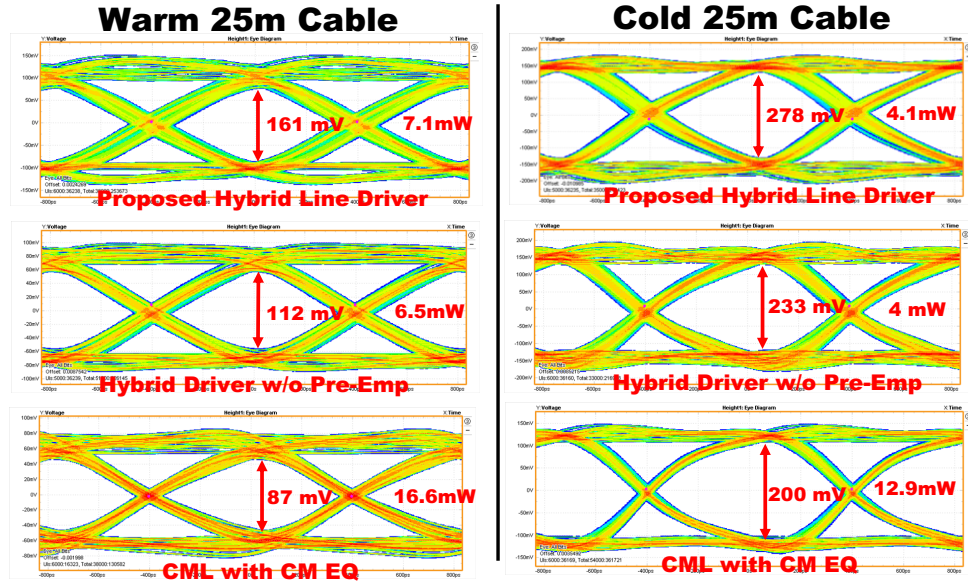


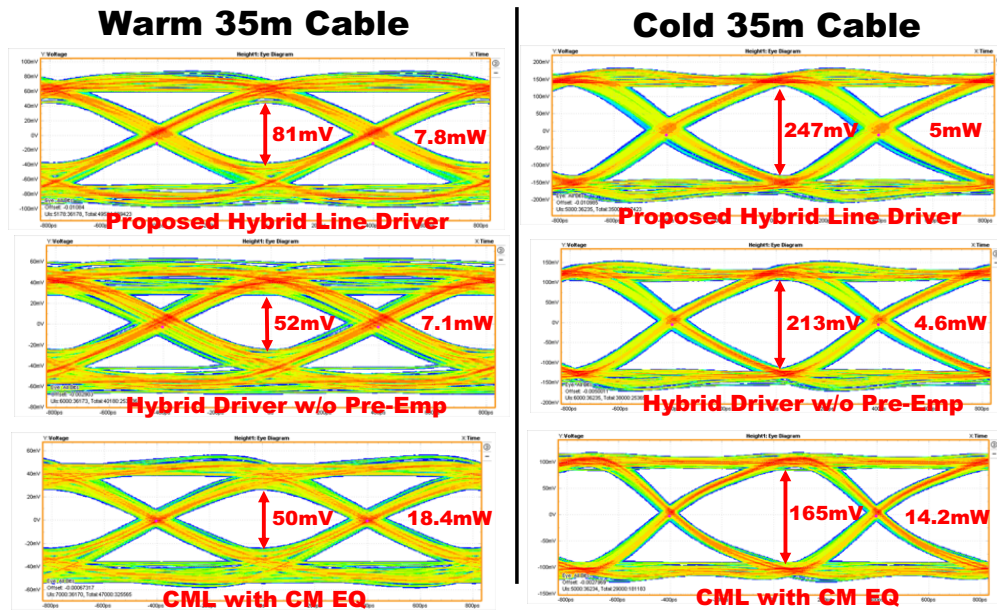
Figure. 43. The measured transmitter output waveforms and eye diagram at the end of 25m twinax cable driven by CML driver without TXEQ and pre-emphasis.

Fig. 43 shows the differential timing waveforms of the transmitted data stream after 25m twinax cable and its totally closed eye diagram. These results are also obtained by programming the driver to a conventional CML driver and turning off the TXEQ and VM cells. Different amplitude for high frequencies and low frequencies results from the large channel loss. Long rise and fall time on data edges are caused by large capacitive load and low channel bandwidth.





(a)



(b)

Figure. 44. Measured eye-diagrams after (a) 25m twinax cable and (b) 35 m twinax cable for PRBS7 data pattern.



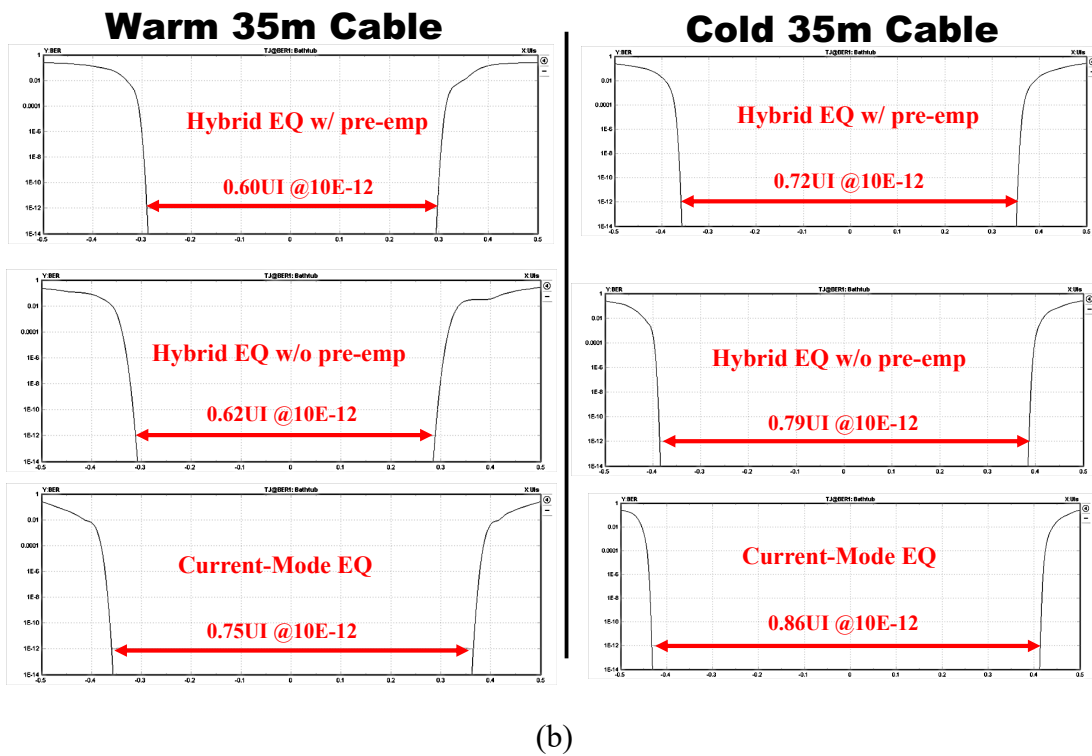
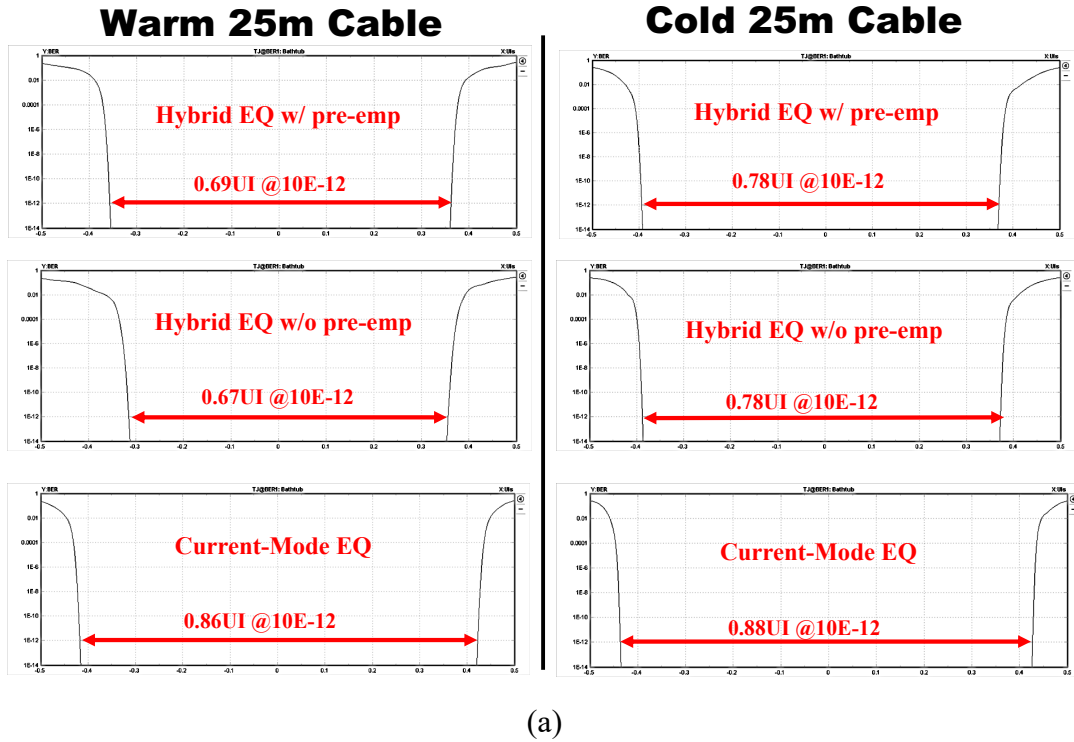


Figure. 45. Measured bathtub curve after (a) 25m twinax cable and (b) 35 m twinax cable for PRBS7 data pattern.

Fig. 44 (a) and (b) show the measured eye diagrams after 25m twinax cable and 35 m twinax cable for PRBS7 data pattern, which is driven by the proposed hybrid line driver with CM TXEQ and VM pre-emphasis. Fig. 45 (a) and (b) show the measured bathtub curves corresponding to the eye diagrams in Fig. 44. The CML driver with only CM TXEQ consumes most of the power among all three configurations but obtains smallest eye open height. Compared to the CML driver, the hybrid drivers saves power consumption significantly and increases the eye height by using the VM SST main driver. The only advantage that the CML driver has over the hybrid ones is better jitter performance (larger eye width) since there is no data-dependent current injected into the supply rails. Although the CML driver obtains a bit better eye width, all of three configurations get widely opened eye diagrams that meet the DUNE BER requirements (less than  $10^{-12}$ ).

Table 1 Simulated comparator noise and power consumption

<b>PRBS7</b>	<b>25m Warm Cable</b>			<b>25m Cold Cable</b>		
	Eye Height (mV)	Eye Width (ps)	Power (mW)	Eye Height (mV)	Eye Width (ps)	Power (mW)
Proposed Hybrid	161	537	7.1	278	613	4.1
Hybrid w/o Pre-Emp	112	523	6.5	233	607	4
CML with CM EQ	87	670	16.6	200	690	12.9

<b>PRBS15</b>	<b>25m Warm Cable</b>			<b>25m Cold Cable</b>		
	Eye Height (mV)	Eye Width (ps)	Power (mW)	Eye Height (mV)	Eye Width (ps)	Power (mW)
Proposed Hybrid	134	360	7.1	243	459	4.1
Hybrid w/o Pre-Emp	76	394	6.5	204	493	4
CML with CM EQ	73	492	16.6	180	628	12.9

<b>PRBS7</b>	<b>35m Warm Cable</b>			<b>35m Cold Cable</b>		
	Eye Height (mV)	Eye Width (ps)	Power (mW)	Eye Height (mV)	Eye Width (ps)	Power (mW)
Proposed Hybrid	81	468	7.8	247	566	5
Hybrid w/o Pre-Emp	52	488	7.1	213	615	4.6
CML with CM EQ	50	582	18.4	165	673	14.2

<b>PRBS15</b>	<b>35m Warm Cable</b>			<b>35m Cold Cable</b>		
	Eye Height (mV)	Eye Width (ps)	Power (mW)	Eye Height (mV)	Eye Width (ps)	Power (mW)
Proposed Hybrid	51	245	7.8	212	452	5
Hybrid w/o Pre-Emp	32	165	7.1	181	505	4.6
CML with CM EQ	45	391	18.4	153	615	14.2

Table 1 extends the measurement results to both PRBS7 and PRBS15 data patterns with the cable length of 25 m and 35 m. Compared to the CML driver with only CM TXEQ, the proposed hybrid line driver consumes 57%-68% less power but obtains 13%-85% more of the eye height. Compared to the hybrid driver without pre-emphasis, the proposed hybrid line driver gets 16%-76% more of the eye height, and only takes 2.5%-9% more power. The BER performance of the proposed hybrid driver is evaluated using the Xilinx KC705 evaluation kit, which is measured using PRBS15 data pattern and warm 25 m cable. Data stream is running for three weeks and a total of  $2.38 \times 10^{15}$  bits are received with no error. Assuming Poisson statistics, this corresponds to a BER of better than  $10^{-15}$ . These results have shown the effectiveness of the proposed hybrid line driver. Given the large eye height and low power consumption, the presented hybrid driver is considered as a more suitable architecture for the DUNE application.

Table 2 shows the comparison of this work with other state of the art with similar data rate. The proposed hybrid line driver compensates the largest channel loss, and achieves comparable energy efficiency with other works.

Table 2 Simulated comparator noise and power consumption

	[3]	[5]	[7]	[11]	[13]	This Work	
<b>Process (nm)</b>	65 nm	90 nm	90 nm	45 nm SOI	90 nm	<b>65 nm</b>	
<b>Supply (V)</b>	1.05 V	1.15 V	1.2 V	1.08 V	1.2 V	<b>1.1 V</b>	
<b>Driver Topology</b>	CM CML	VM	VM	VM SST	Hybrid	<b>Hybrid</b>	
<b>Data Rate (Gbps)</b>	15 Gbps	4 Gbps	5 Gbps	7.4 Gbps	6 Gbps	<b>1.28 Gbps</b>	
<b>TX Swing (Vpp)</b>	160 mVpp	500 mVpp	400 mVpp	800 mVpp	300 mVpp	<b>1.1 Vpp</b>	
<b>Channle loss at Nyquist (dB)</b>	10 dB	8-12 dB	6 dB	2.6 dB	3.7 dB	<b>25 dB</b>	<b>13 dB</b>
<b>Power (mW)</b>	34 mW	17 mW	4.9 mW	32 mW	4.1 mW	<b>7.8 mW</b>	<b>4.1 mW</b>
<b>Energy Efficiency (pJ/bit)</b>	2.3 pJ/b	4.25 pJ/b	0.98 pJ/b	4.3 pJ/b	1.26 pJ/b	<b>6.1 pJ/b</b>	<b>3.1 pJ/b</b>
<b>TX Area (mm^2)</b>	0.033	0.02	0.023	N/A	0.035	<b>0.023</b>	

## CHAPTER 5 CONCLUSION

### 5.1 Summary

A new asynchronous serial transceiver that supports an auxiliary channel yielding additional data transmission capability is demonstrated in a  $0.13\text{mm}^2$  65nm CMOS IC. The prototype transceiver allows for both its primary and auxiliary data streams, at 2.56Gbps and 80Mbps respectively, to be recovered simultaneously with good jitter and BER performance. The analysis of the auxiliary data rate that can be achieved from available channel bandwidth margin is given in this dissertation. The contribution supports secure methods by offering a way to utilize an additional auxiliary channel and extra data bandwidth to support potential security measures such as the inclusion of authentication data, additional support of encryption or other methods requiring another channel or more bandwidth, steganography, etc. The additional novelty is that this auxiliary channel in serial transceiver are provided in a way that offers backward compatibility, interoperability with non-equipped designs, and minimal redesign of existing systems.

A hybrid transmitter with a proposed line driver equips VM pre-emphasis and CM TXEQ has been designed, implemented and measured at cryogenic temperature to compensate large frequency-dependent channel loss introduced by 25-35 meter twinax cables. Large eye height is obtained by employing SST-based VM main driver and VM pre-emphasis together with the CM TXEQ architecture. The combination of the VM SST cells and CM TXEQ also offers the benefits of low power consumption and low design complexity.

## **5.2 Future Work**

The further development on the dissertation topics is to boost the transmitted data rate. The goal of the first work in this dissertation is to propose a new modulation and demodulation scheme based on the conventional transceiver architecture to provide an auxiliary channel and extra bandwidth. This work is our first prototype, and its purpose is to verify the validity of the new scheme. Thus this prototype is not operating at a very high speed or over a high-loss channel, and there is no equalization technique being used in the first work. In the future, boosting the data rate and maintaining low power consumption in the transceiver should be given priority attention in the future study.

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